

Roadmap

- **What is a Computer Network?**
- Applications of Networking
- Classification of Networks
- Layered Architecture
- Network Core
- Delay & loss in packet-switched networks
- Internet Structure
- Transmission Media (self study)
- History (self study)

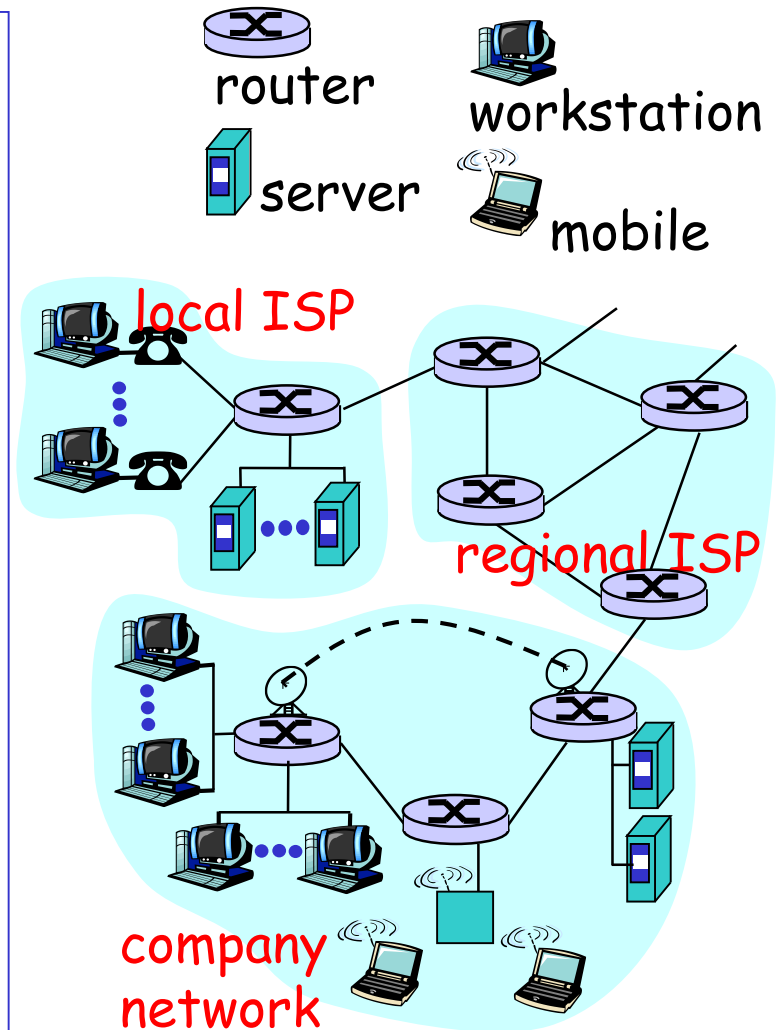
Computer Network?



- ❑ “interconnected collection of autonomous computers connected by a *single* technology” [Tanenbaum]
- ❑ What is the Internet?
 - “network of networks”
 - “collection of networks interconnected by routers”
 - “a communication medium used by millions”
 - Email, chat, Web “surfing”, streaming media
- ❑ Internet ≠ Web

The "nuts and bolts" view of the Internet

- ❑ Network Edge
- ❑ millions of connected computing devices: *hosts, end-systems*
 - PCs workstations, servers
 - PDAs phones, toastersrunning *network apps*
- ❑ *communication links*
 - fiber, copper, radio, satellite
 - Links have different *bandwidth*
- ❑ *routers*: forward packets
- ❑ *Packet*: a piece of messg.



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Applications (1)

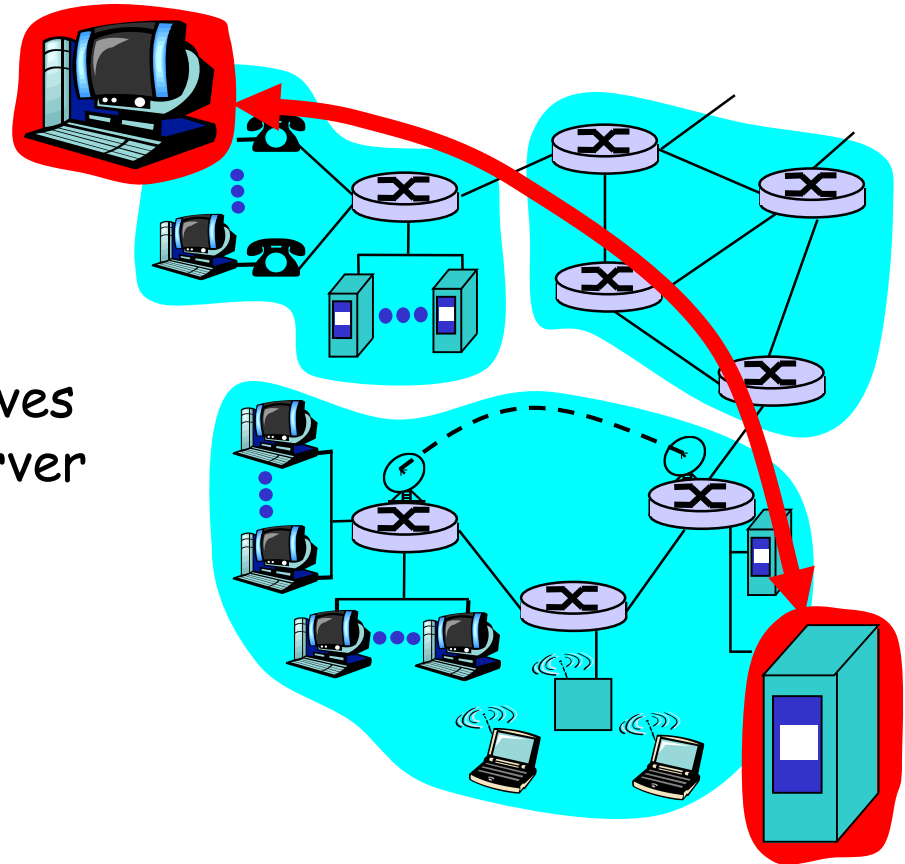
□ end systems (hosts):

- run application programs
- e.g. Web, email
- at "edge of network"

□ client/server model

- client host requests, receives service from always-on server
- e.g. Web browser/server; email client/server

□ Client/server model is applicable in an *intranet*.

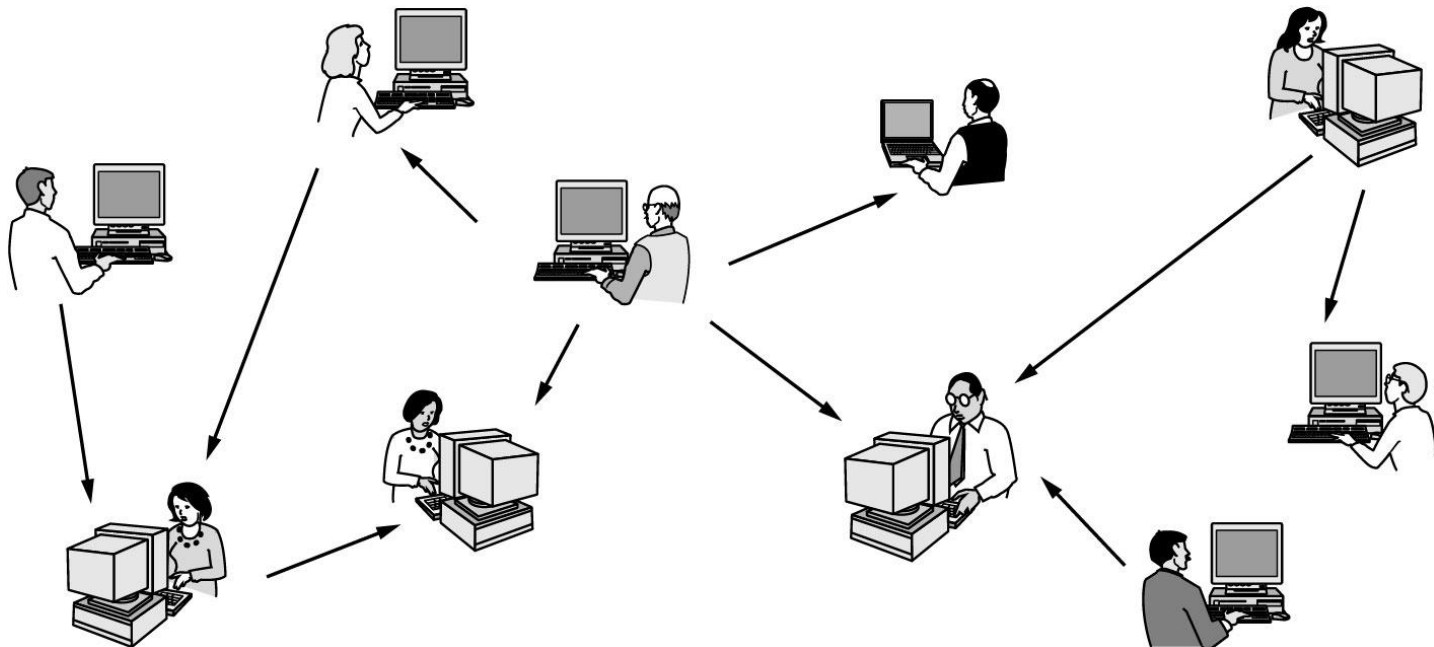


Applications (2)

□ peer-peer model:

- No fixed clients or servers
- Each host can act as both client & server

□ Examples: Napster, Gnutella, KaZaA



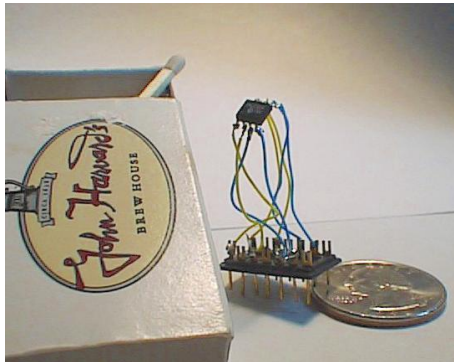
Applications (3)

- ❑ WWW
- ❑ Instant Messaging (Internet chat, text messaging on cellular phones)
- ❑ Peer-to-Peer
- ❑ Internet Phone
- ❑ Video-on-demand
- ❑ Distributed Games
- ❑ Remote Login (SSH client, Telnet)
- ❑ File Transfer

"Cool" Appliances



IP picture frame
<http://www.ceiva.com/>



World's smallest web server
<http://www-ccs.cs.umass.edu/~shri/iPic.html>



Web-enabled toaster+weather forecaster

Roadmap

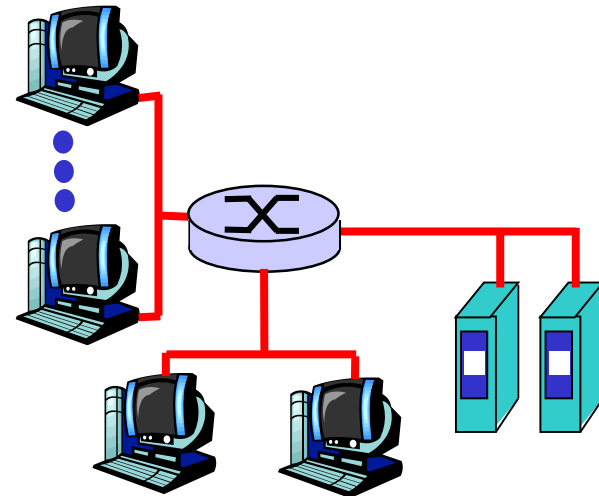
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A Classification of Networks

- ❑ Local Area Network (LAN)
- ❑ Metropolitan Area Network (MAN)
- ❑ Wide Area Network (WAN)
- ❑ Wireless LANs & WANs
- ❑ Home Networks

Local Area Network (LAN)

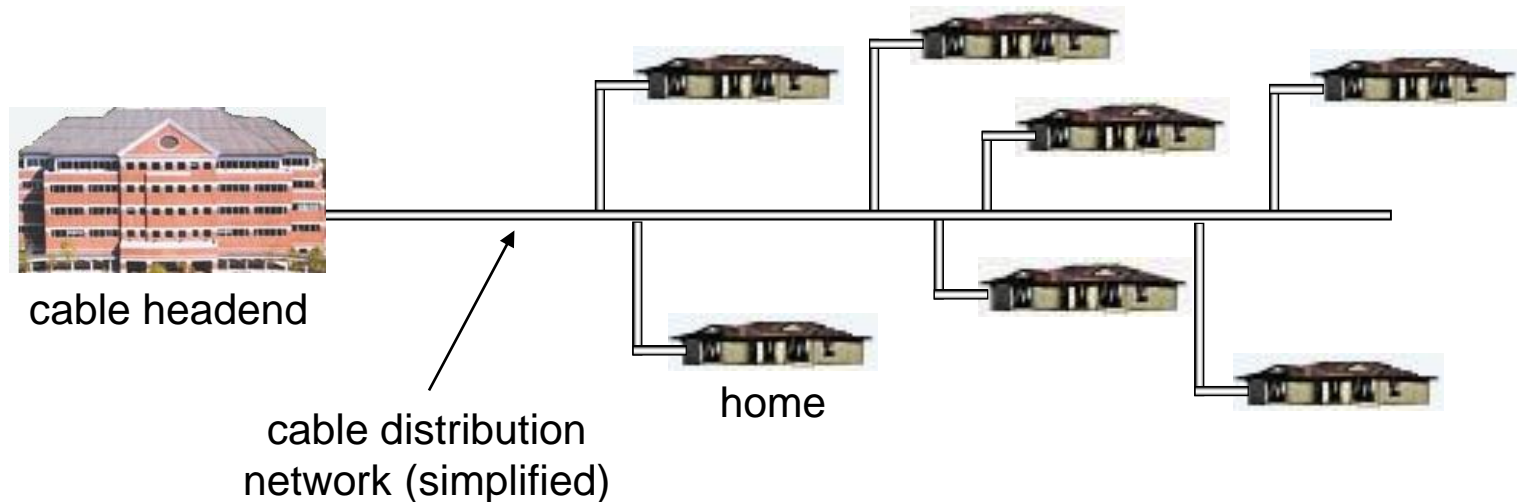
- ❑ company/univ **local area network** (LAN) connects end system to edge router
- ❑ **Ethernet:**
 - shared or dedicated link connects end system and router
 - 10 Mbs, 100Mbps, Gigabit Ethernet
- ❑ **deployment:** institutions, home LANs happening now



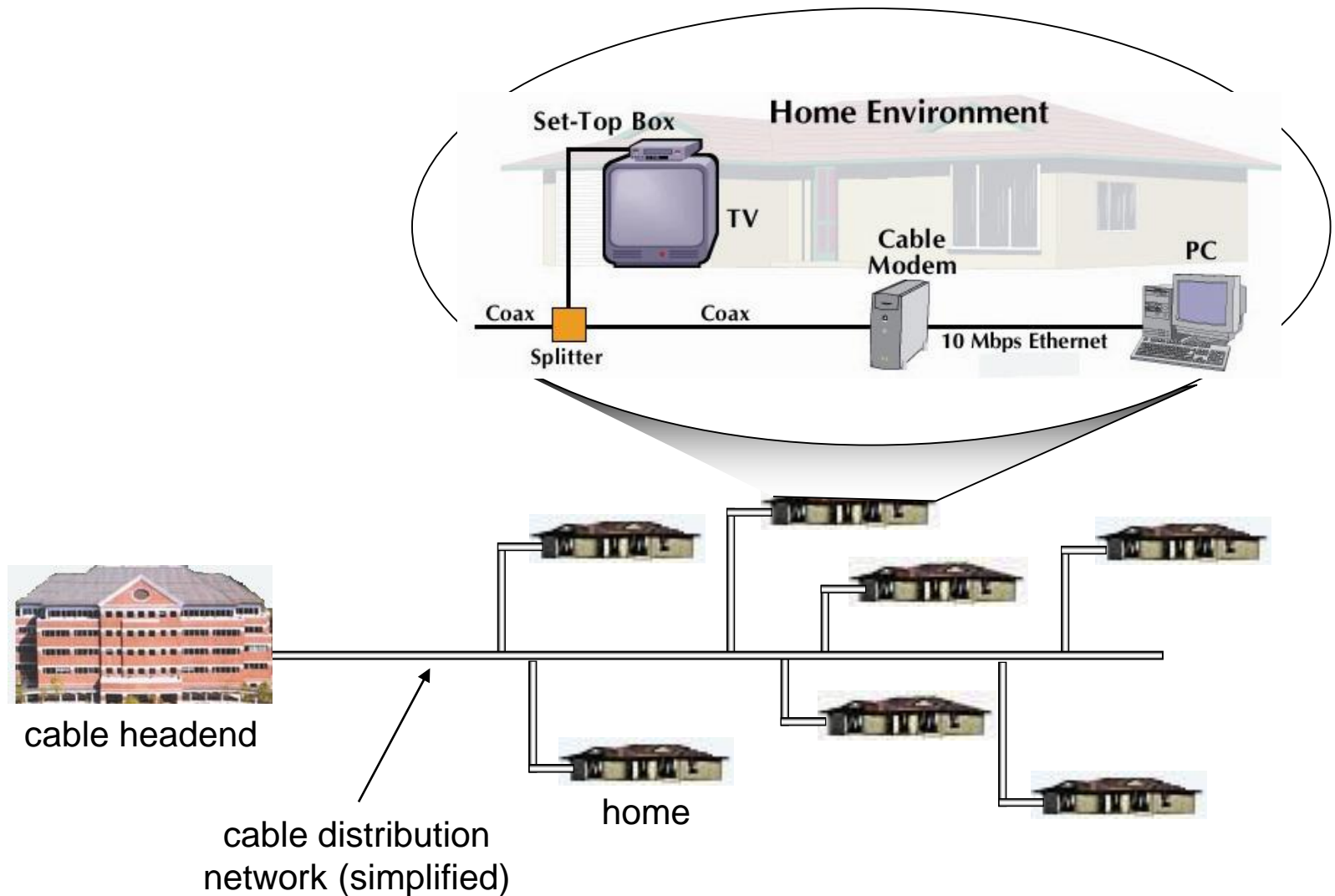
Metropolitan Area Network (MAN)

A Cable TV Network is an example of a MAN

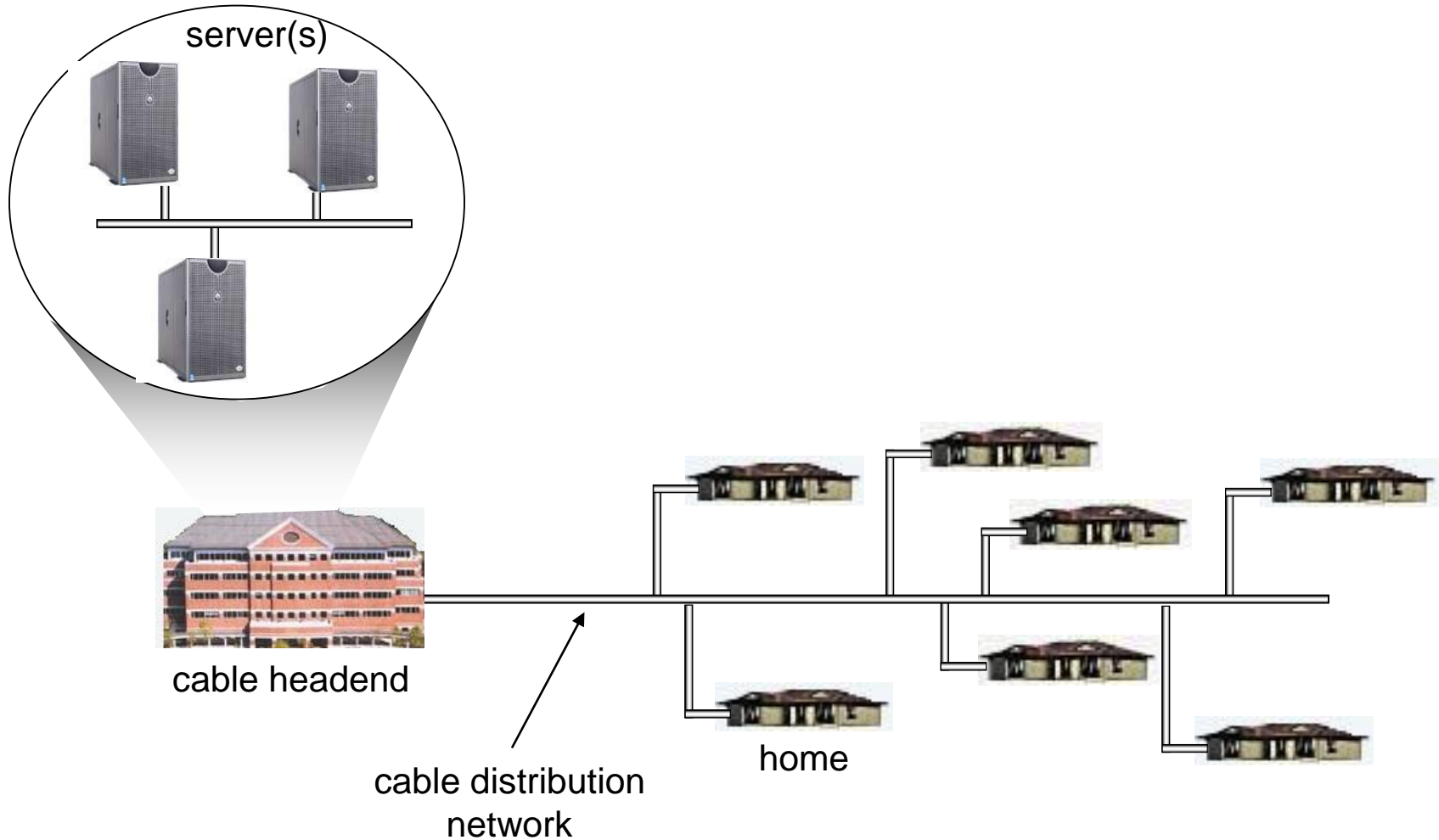
Typically 500 to 5,000 homes



Cable Network Architecture: Overview

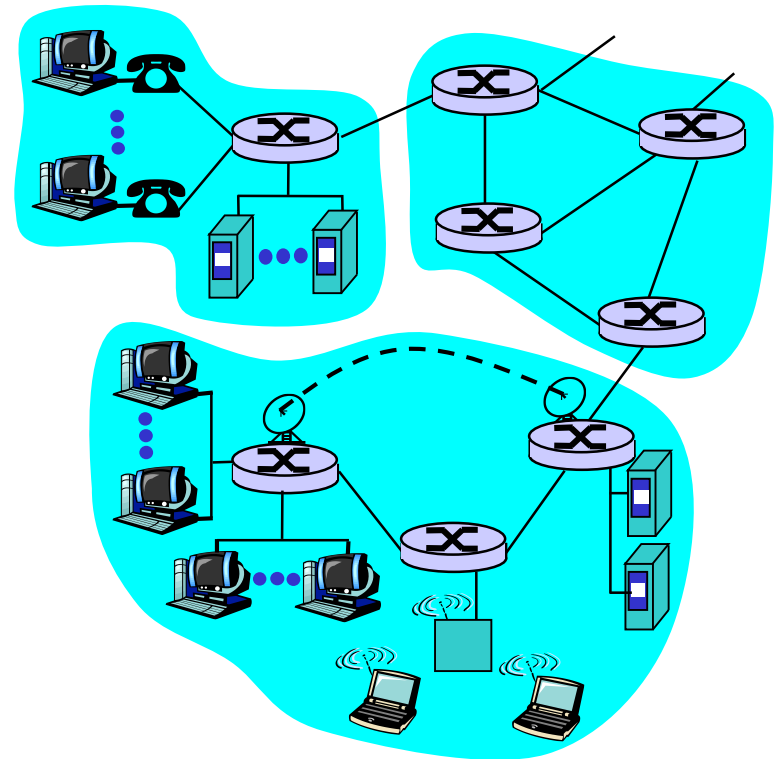


Cable Network Architecture: Overview



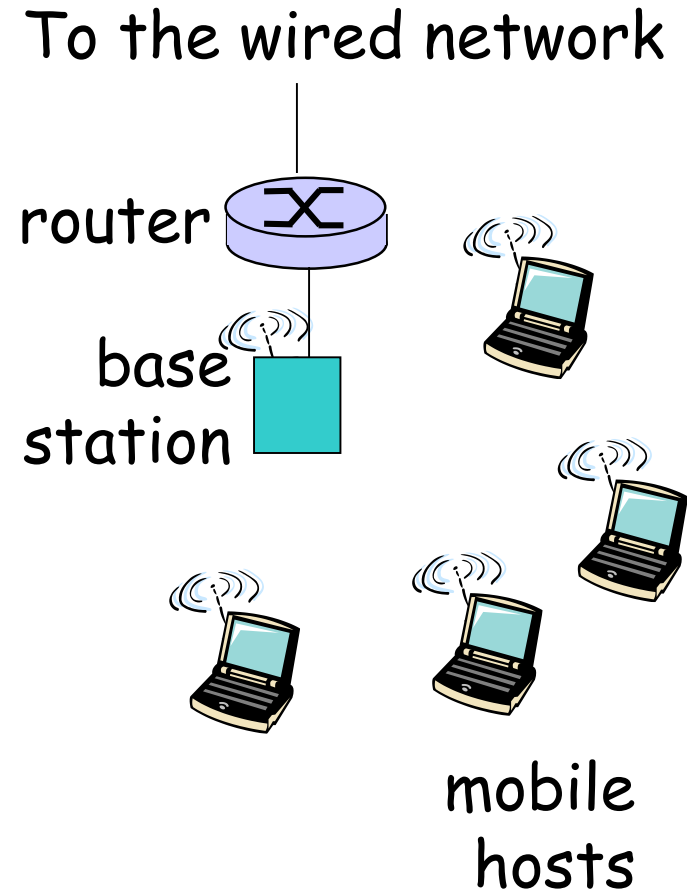
Wide Area Network (WAN)

- ❑ Spans a large geographic area, e.g., a country or a continent
- ❑ A WAN consists of several transmission lines and routers
- ❑ Internet is an example of a WAN



Wireless Networks

- shared *wireless* access network connects end system to router
 - via base station aka "access point"
- **wireless LANs:**
 - 802.11b (WiFi): 11 Mbps
- **wider-area wireless access**
 - provided by telco operator
 - 3G ~ 500 kbps - 3MB
 - 4G 3MB-300 MB
 - WAP/GPRS in Europe

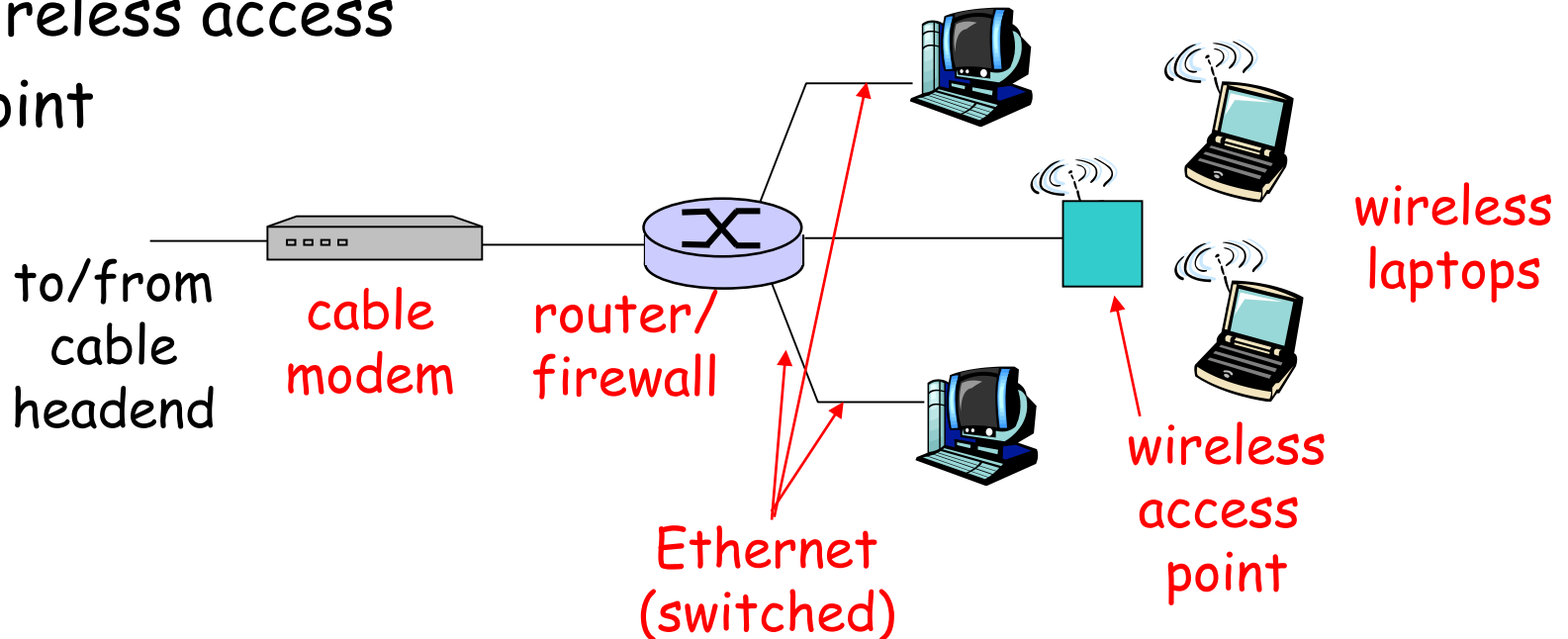




Home networks

Typical home network components:

- ❑ ADSL or cable modem
- ❑ router/firewall/NAT
- ❑ Ethernet
- ❑ wireless access point



"internetworking"?

- ❑ internetwork - interconnection of networks - also called an "internet"
- ❑ Subnetwork - a constituent of an internet
- ❑ Intermediate system - a device used to connect two networks allowing hosts of the networks to correspond with each other
 - Bridge
 - Routers
- ❑ **Internet is an example of an internetwork.**

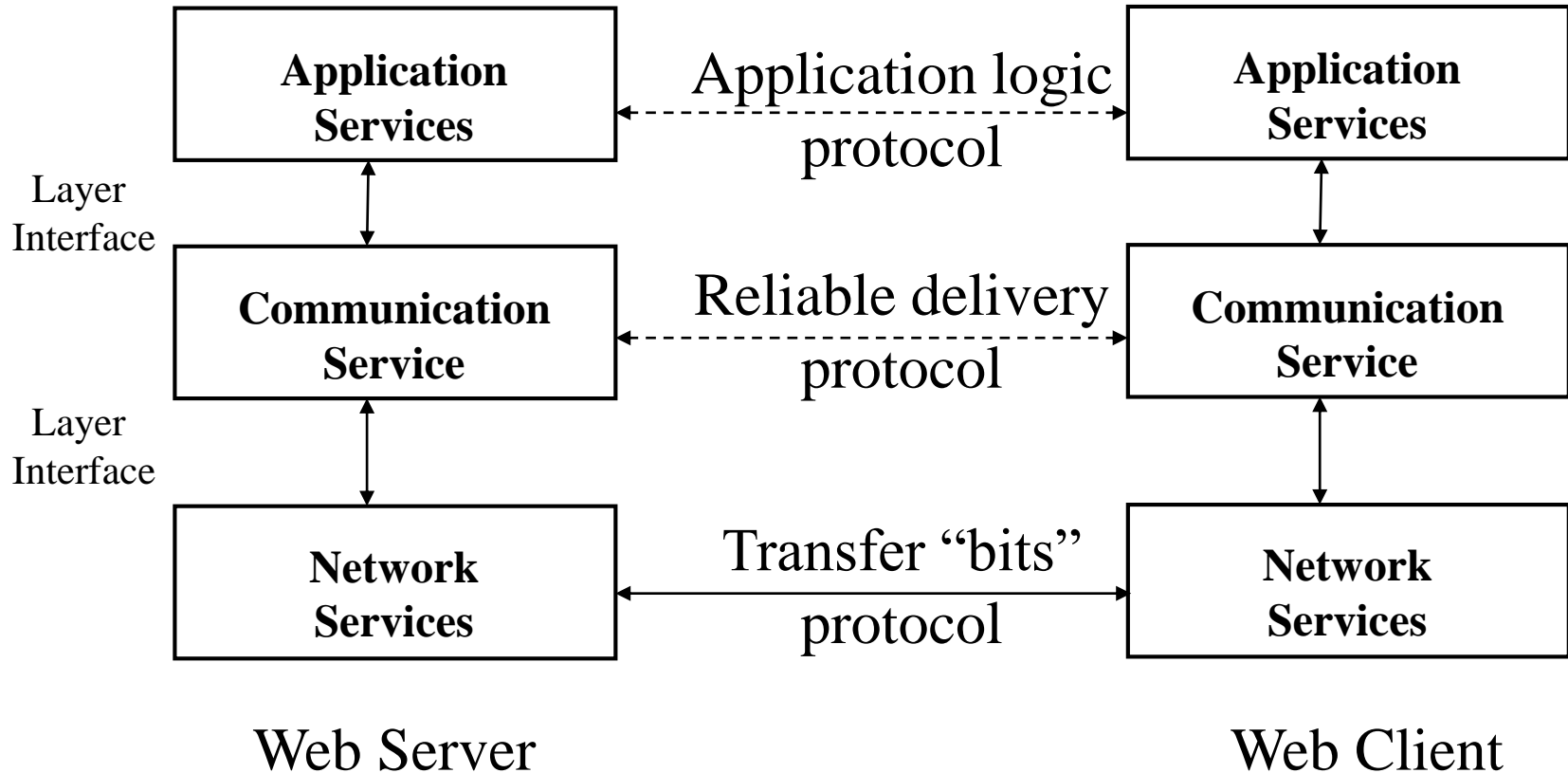
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Layered Architecture: Why?

- ❑ Networks are complex with many pieces
 - Hosts, routers, links, applications, protocols, hardware, software
- ❑ Can we organize it, somehow?
- ❑ Let's consider a Web page request:
 - Browser requests Web page from server
 - Server should determine if access is privileged
 - Reliable transfer page from server to client
 - Physical transfer of "bits" from server to client

Layers, Protocols, Interfaces



Layered Architecture (Review 1/2)

- ❑ Networks organized as a stack of layers?
 - The purpose of a layer is to offer services to the layer above it using an interface (programming language analogy: libraries hide details while providing a service)
 - Reduces design complexity
- ❑ Protocols: peer-to-peer layer-n conversations
- ❑ Data Transfer: each layer passes data & control information to the layer below; eventually physical medium is reached.

Review (2/2)

- A set of layers & protocols is called a Network Architecture. These specifications enable hardware/software developers to build systems compliant with a particular architecture.
 - E.g., TCP/IP, OSI

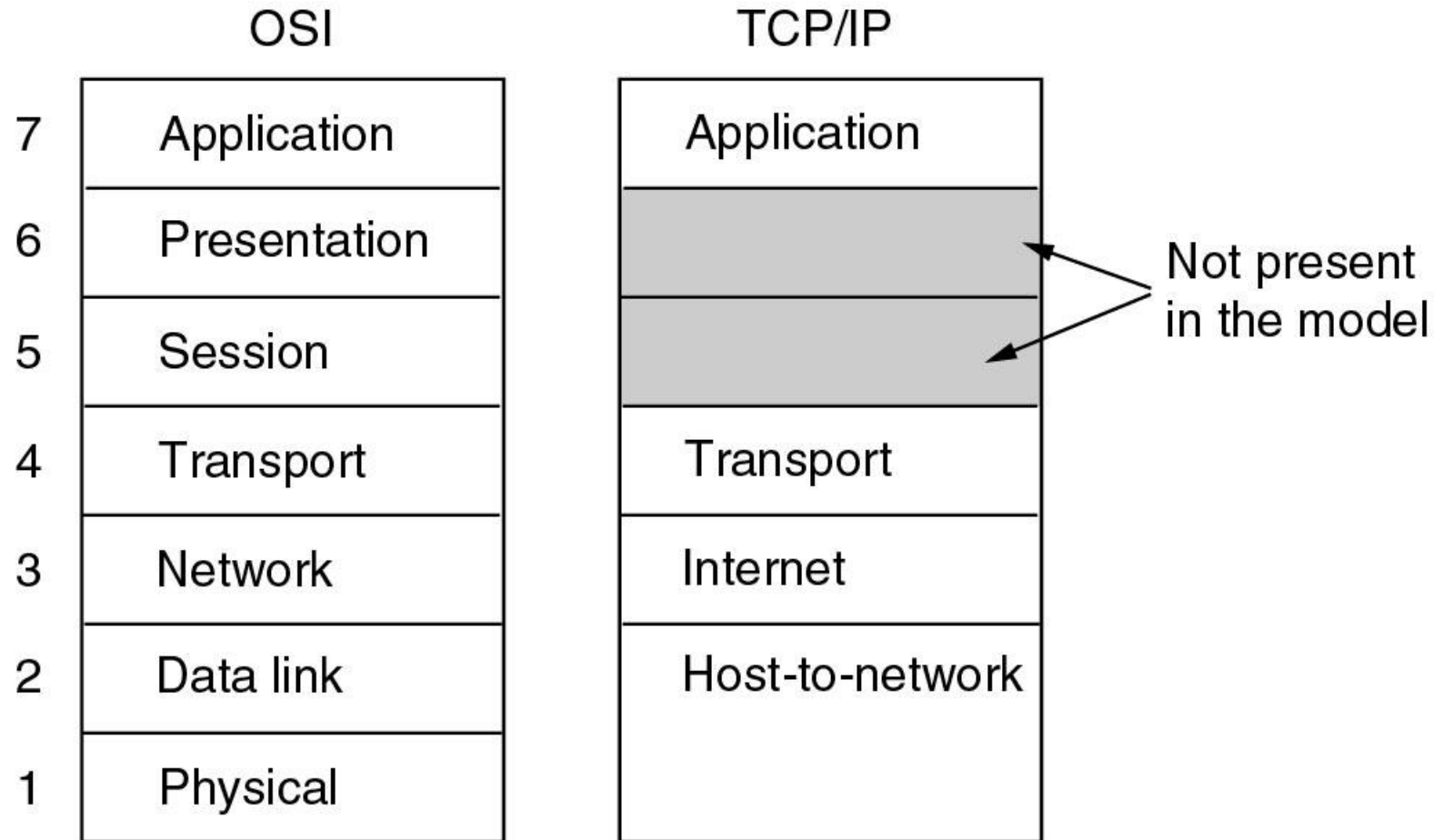
Layering: Design Issues

- ❑ Identify senders/receivers?
 - Addressing
- ❑ Unreliable physical communication medium?
 - Error detection
 - Error control
 - Message reordering
- ❑ Sender can swamp the receiver?
 - Flow control
- ❑ Multiplexing/Demultiplexing

Reference Model

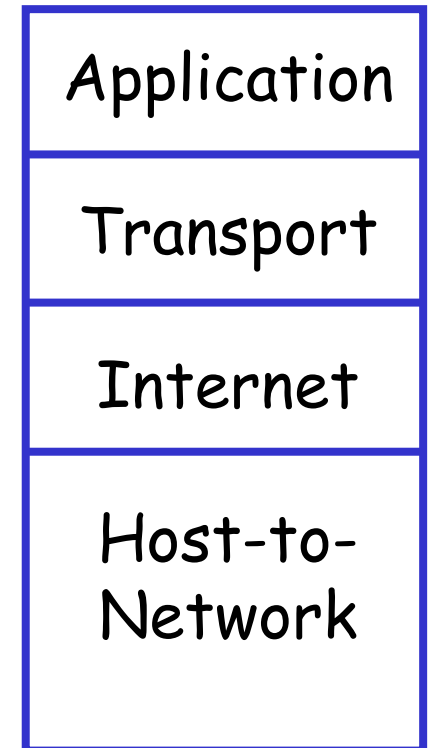
- ❑ Open Systems Interconnection (OSI) Model
- ❑ TCP/IP Model

Reference Models (2)



TCP/IP Model: History

- ❑ Originally used in the ARPANET
- ❑ ARPANET required networks using leased telephone lines & radio/satellite networks to interoperate
- ❑ Goals of the model are:
 - Seamless interoperability
 - Wide-ranging applications
 - Fault-tolerant to some extent



The Internet Layer

- ❑ End systems inject datagrams in the networks
- ❑ A transmission path is determined for each packet (routing)
- ❑ A "best effort" service
 - Datagrams might be lost
 - Datagrams might be arrive out of order
- ❑ Analogy: Postal system

The Transport Layer

- ❑ Concerned with end-to-end data transfer between end systems (hosts)
- ❑ Transmission unit is called segment
- ❑ TCP/IP networks such as the Internet provides two types of services to applications
 - "connection-oriented" service - Transmission Control Protocol (TCP)
 - "connectionless" service - User Datagram Protocol (UDP)

TCP: Connection-oriented Service

- ❑ Handshaking between client & server programs
 - Parameters for ensuing exchange
 - Maintain connection-state
- ❑ Packet switches maintain any connection-state;
 - hence "connection-oriented"
- ❑ Similar to a phone conversation
- ❑ TCP is bundled with reliability, congestion control, and flow control.

UDP: Connectionless Service

- ❑ No handshaking
- ❑ Send whenever and however you want
- ❑ A “best effort” service
 - No reliability
 - No congestion & flow control services
- ❑ Why is it needed?

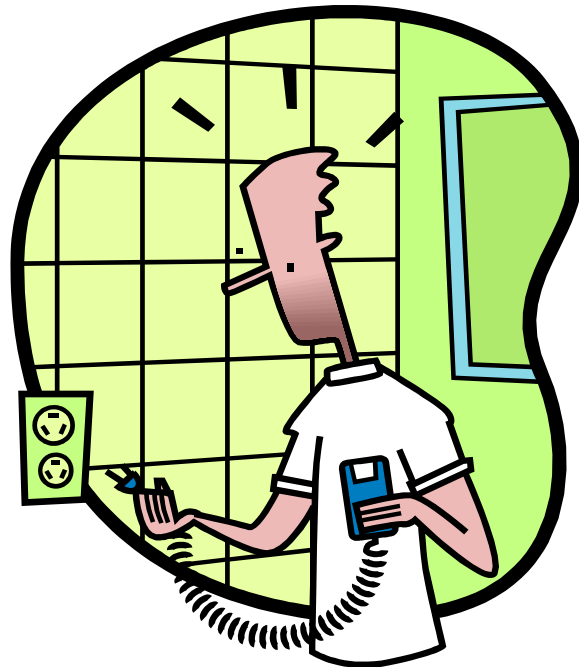


The Application Layer

- ❑ Residence of network applications and their application control logic
- ❑ Examples include:
 - HTTP
 - FTP
 - Telnet
 - SMTP
 - DNS

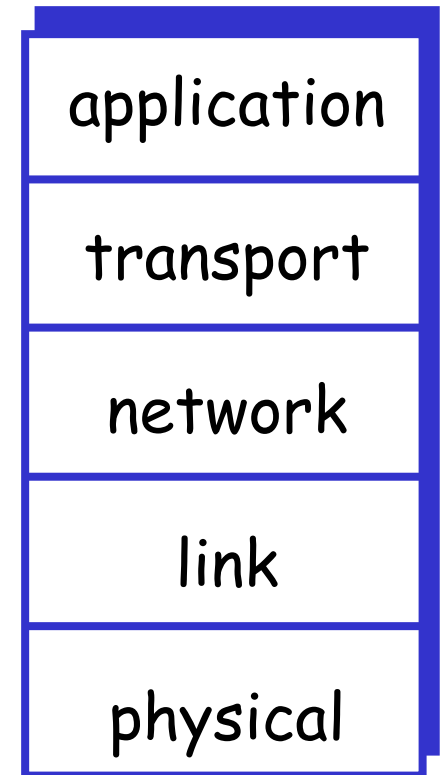
The Host-to-Network Layer

- ❑ Somehow, host has to connect to the network and be able to send IP Datagrams
- ❑ How?



Internet protocol stack

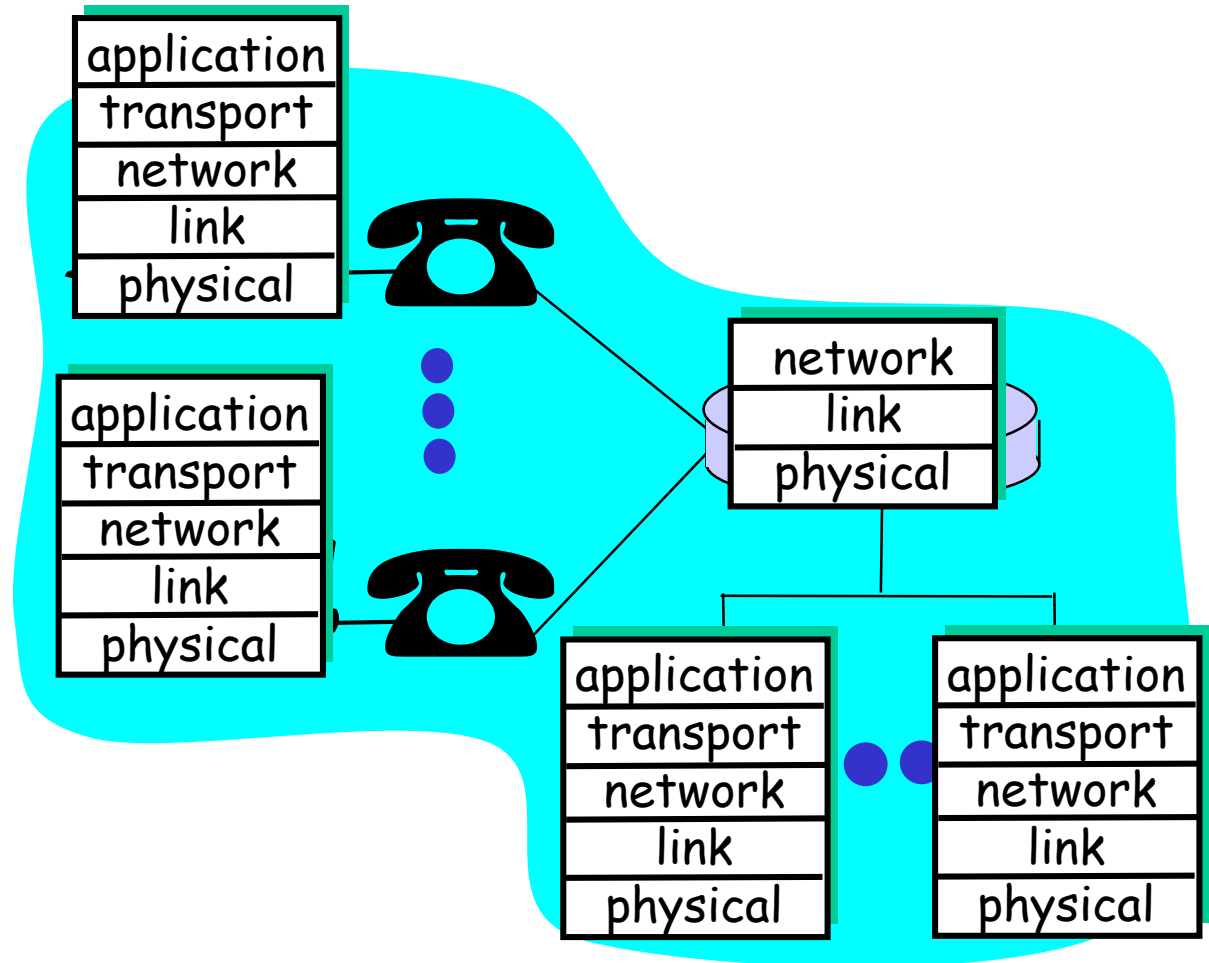
- ❑ **application:** supporting network applications
 - FTP, SMTP, STTP
- ❑ **transport:** host-host data transfer
 - TCP, UDP
- ❑ **network:** routing of datagrams from source to destination
 - IP, routing protocols
- ❑ **link:** data transfer between neighboring network elements
 - PPP, Ethernet
- ❑ **physical:** bits "on the wire"



Layering: logical communication

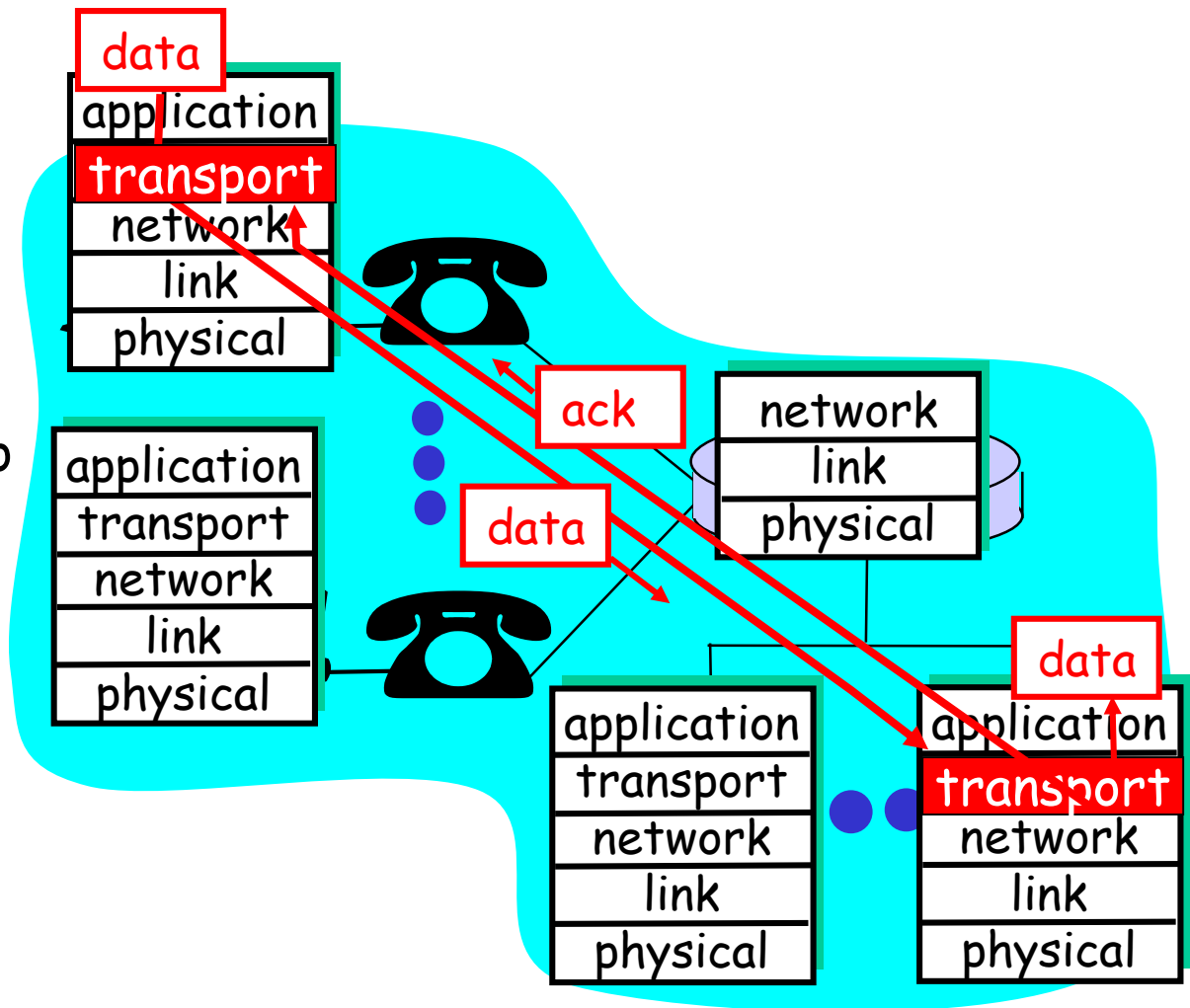
Each layer:

- distributed
- "entities" implement layer functions at each node
- entities perform actions, exchange messages with peers

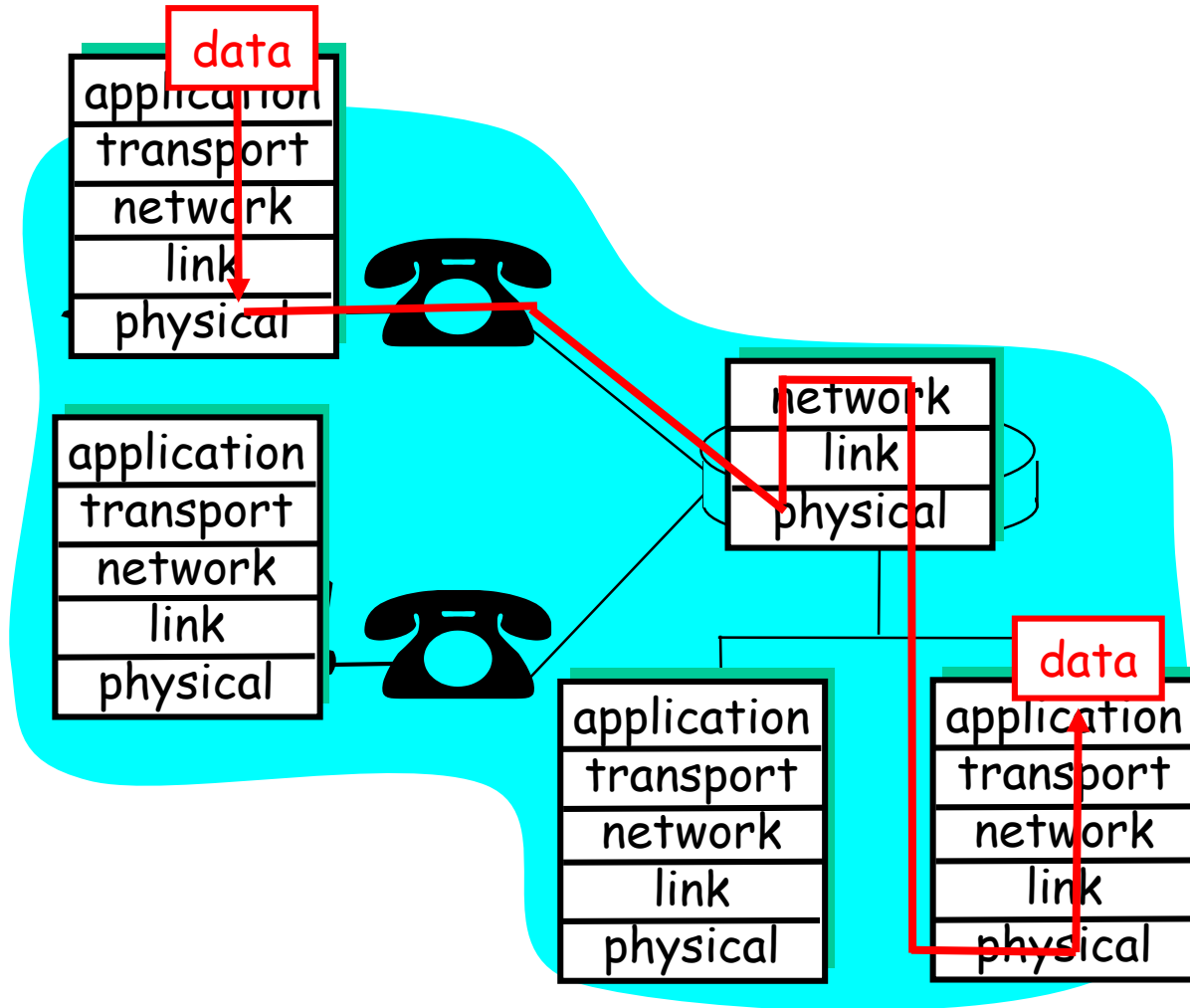


Layering: logical communication

- ❑ take data from app
- ❑ generate "segment" according to transport protocol
- ❑ add addressing, reliability check info to form "datagram"
- ❑ send datagram to peer
- ❑ wait for peer to ack receipt



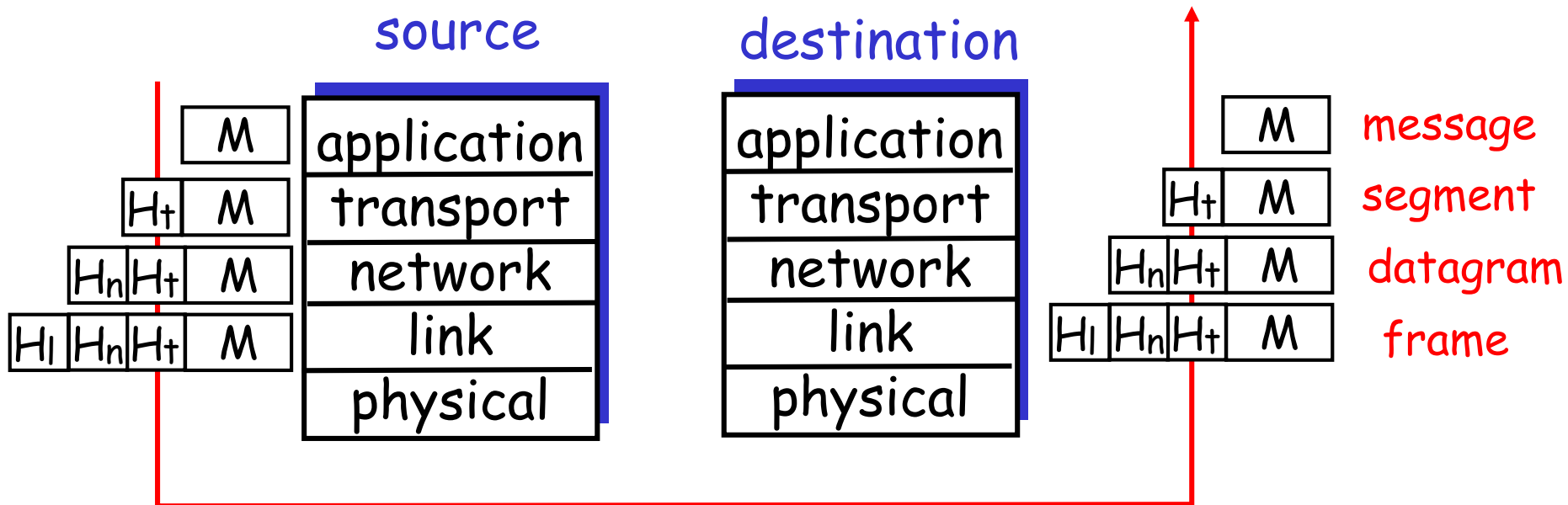
Layering: physical communication



Protocol layering and data

Each layer takes data from above

- adds header information to create new data unit
- passes new data unit to layer below

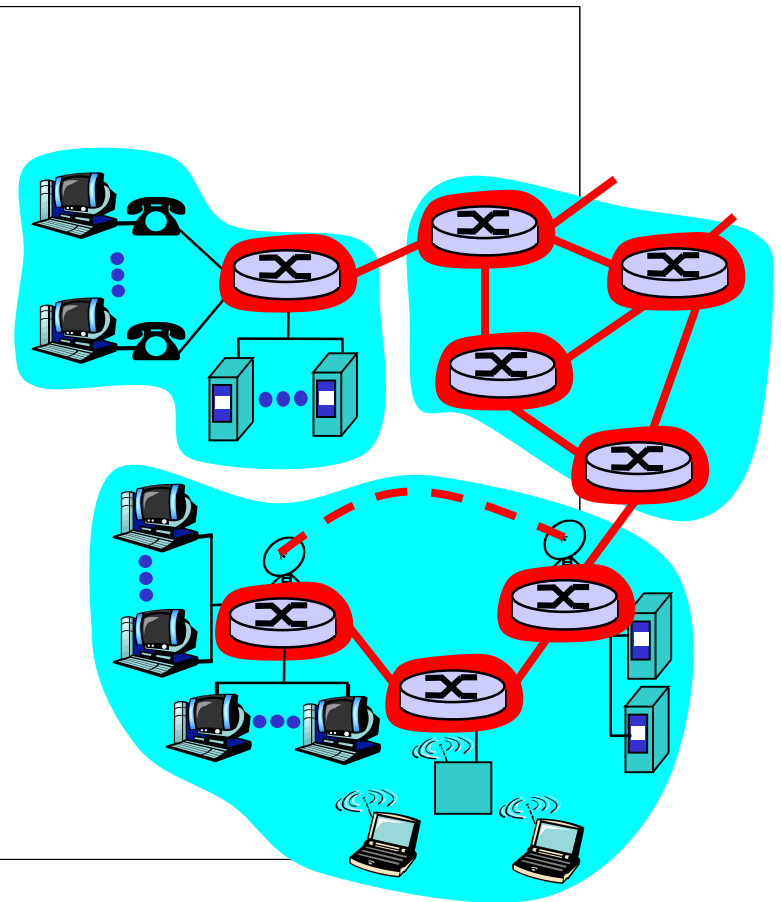


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The Network Core

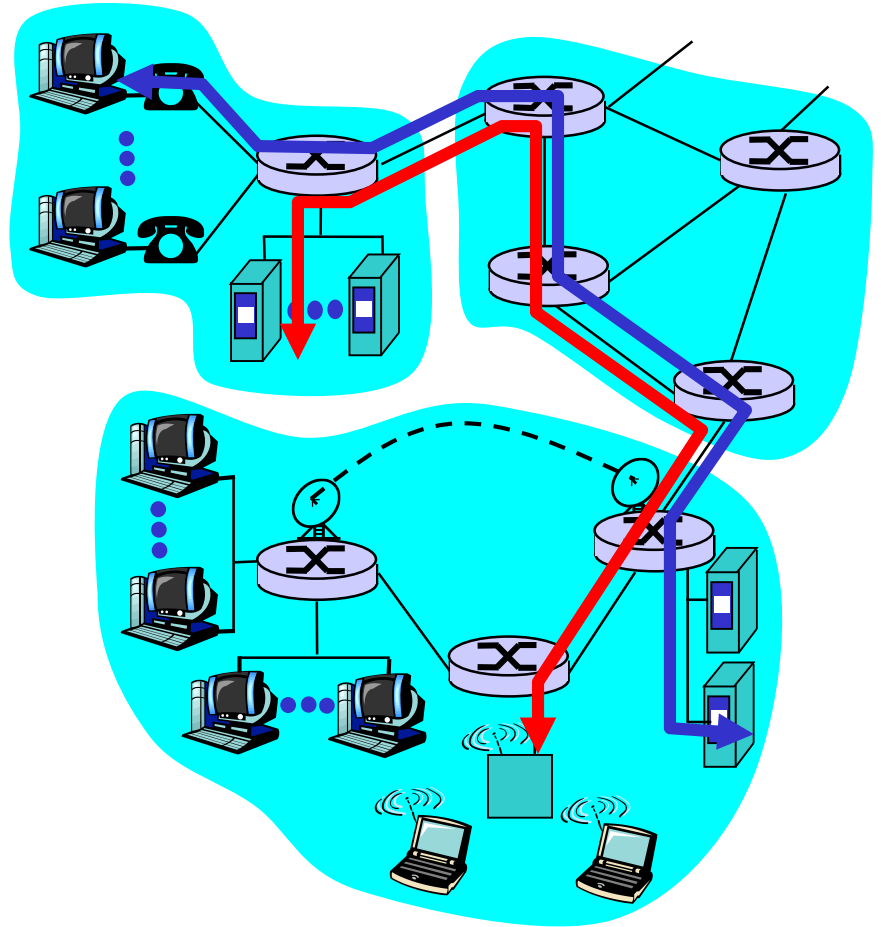
- ❑ mesh of interconnected routers
- ❑ the fundamental question: how is data transferred through net?
 - circuit switching: dedicated circuit per call: telephone net
 - packet-switching: data sent thru net in discrete "chunks"



Network Core: Circuit Switching

End-to-end resources reserved for "call"

- ❑ Link bandwidth, switch capacity
- ❑ Dedicated resources with no sharing
- ❑ Guaranteed transmission capacity
- ❑ Call setup required
- ❑ "Blocking" may occur



Network Core: Circuit Switching

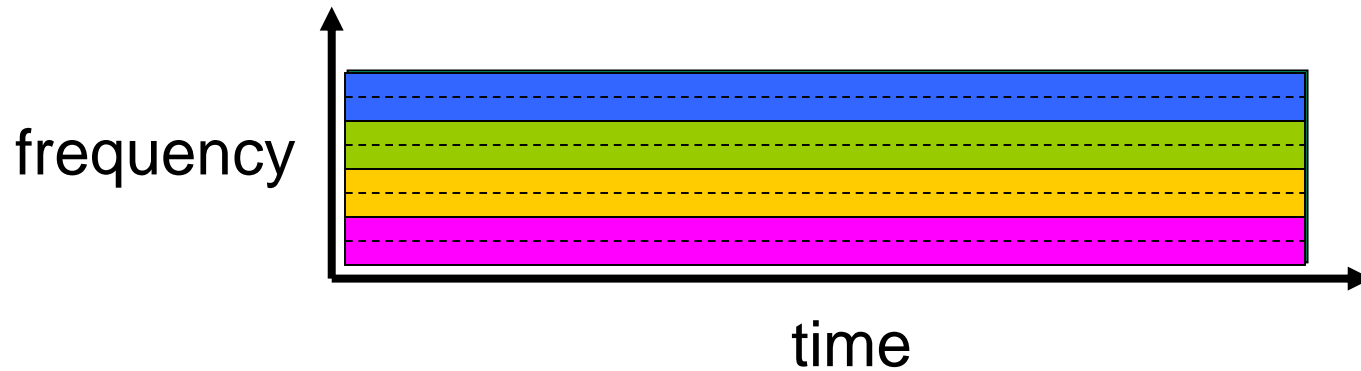
- ❑ Capacity of medium exceeds the capacity required for transmission of a single signal
 - How can we improve "efficiency"? Let's **multiplex**.
- ❑ Divide link bandwidth into "pieces":
 - frequency division - FDMA
 - time division - TDMA

Circuit Switching: FDMA and TDMA

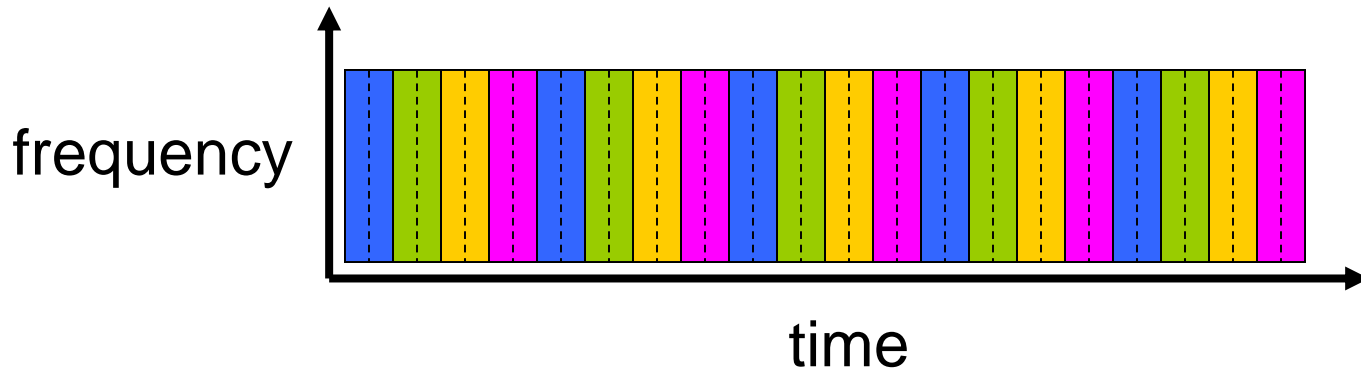
FDMA

Example:

4 users



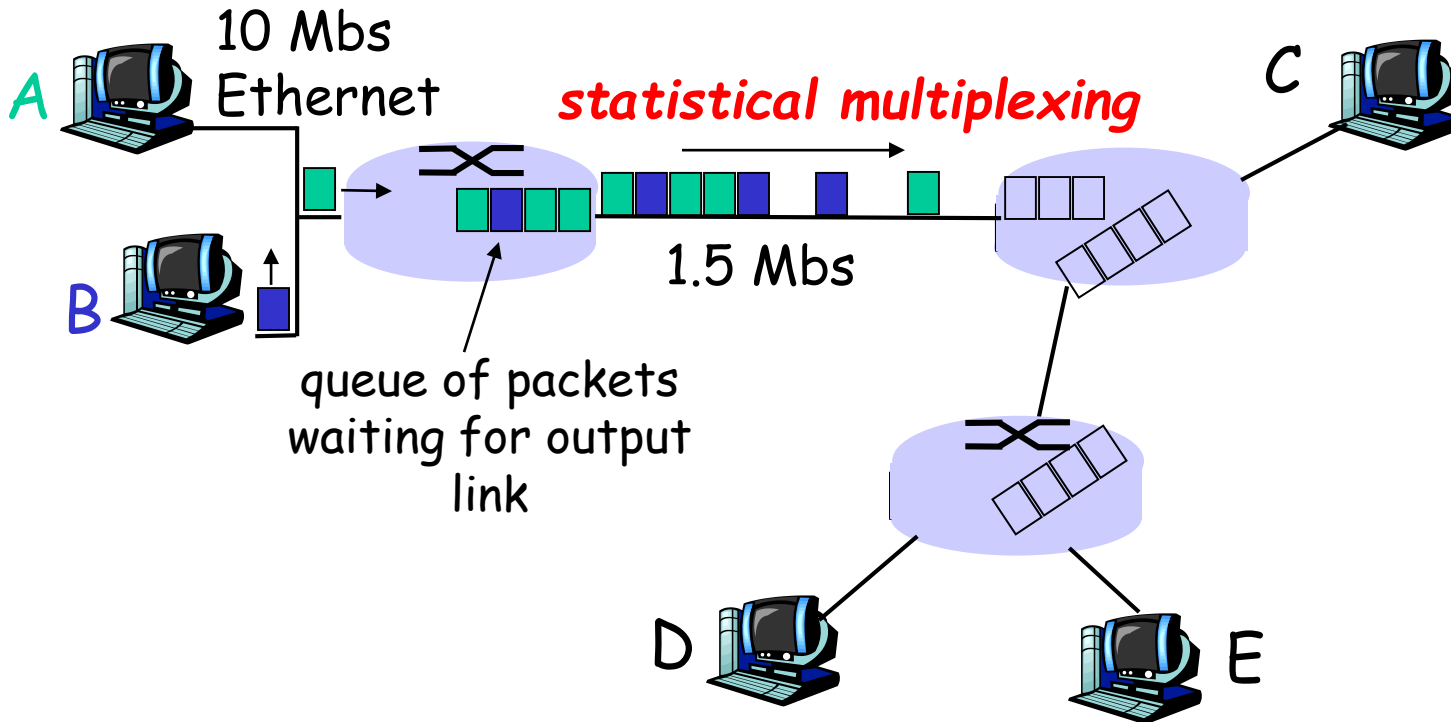
TDMA



Network Core: Packet Switching

- ❑ “store-and-forward” transmission
- ❑ source breaks long messages into smaller “packets”
- ❑ packets *share* network resources
- ❑ each packet uses full link bandwidth
- ❑ resource contention
 - aggregate resource demand can exceed amount available
 - congestion: packets queue, wait for link use

Packet Switching: Statistical Multiplexing



Sequence of A & B packets does not have fixed pattern → **statistical multiplexing**.

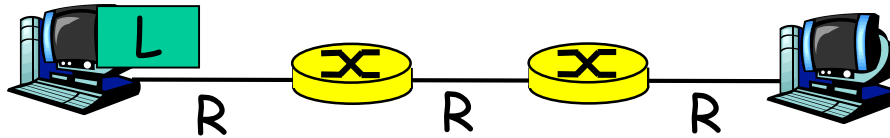
In TDM each host gets same slot in revolving TDM frame.

Packet switching versus circuit switching

Is packet switching a "slam dunk winner?"

- ❑ Great for bursty data
 - resource sharing
- ❑ **Excessive congestion:** packet delay and loss
 - protocols needed for reliable data transfer, congestion control
- ❑ **Q: How to provide circuit-like behavior?**
 - bandwidth guarantees needed for audio/video apps
 - still an unsolved problem

Packet-switching: store-and-forward

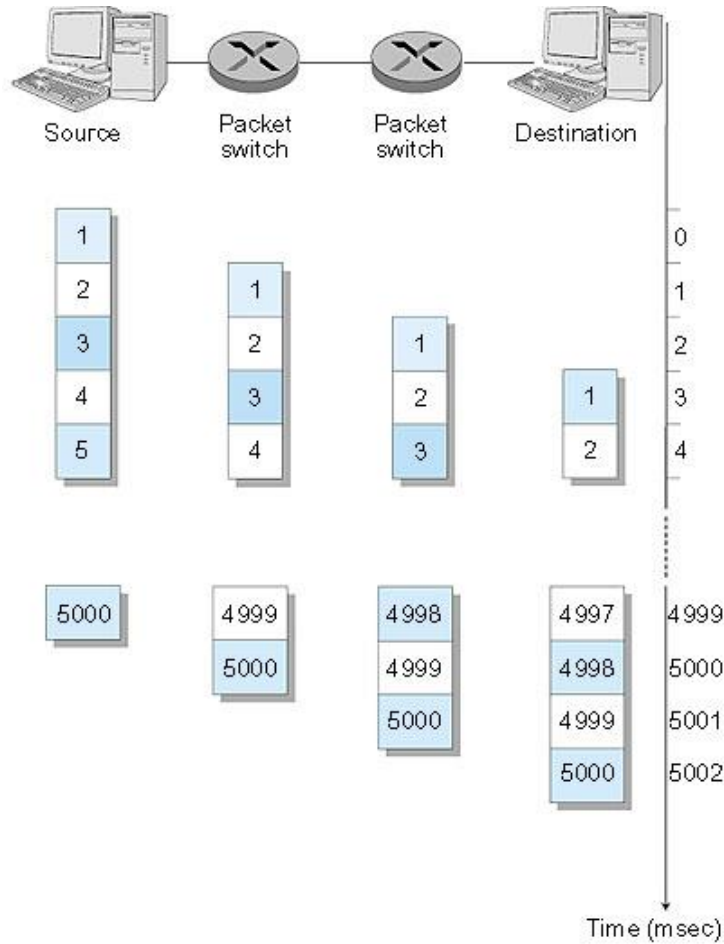


- ❑ Takes L/R seconds to transmit (push out) packet of L bits on to link or R bps
- ❑ Entire packet must arrive at router before it can be transmitted on next link: *store and forward*
- ❑ delay = $3L/R$

Example:

- ❑ $L = 7.5$ Mbits
- ❑ $R = 1.5$ Mbps
- ❑ delay = 15 sec

Packet Switching: Message Segmenting



Now break up the message into 5000 packets

- ❑ Each packet 1,500 bits
- ❑ 1 msec to transmit packet on one link
- ❑ *pipelining*: each link works in parallel
- ❑ Delay reduced from 15 sec to 5.002 sec

Packet-switched networks: forwarding

□ datagram network:

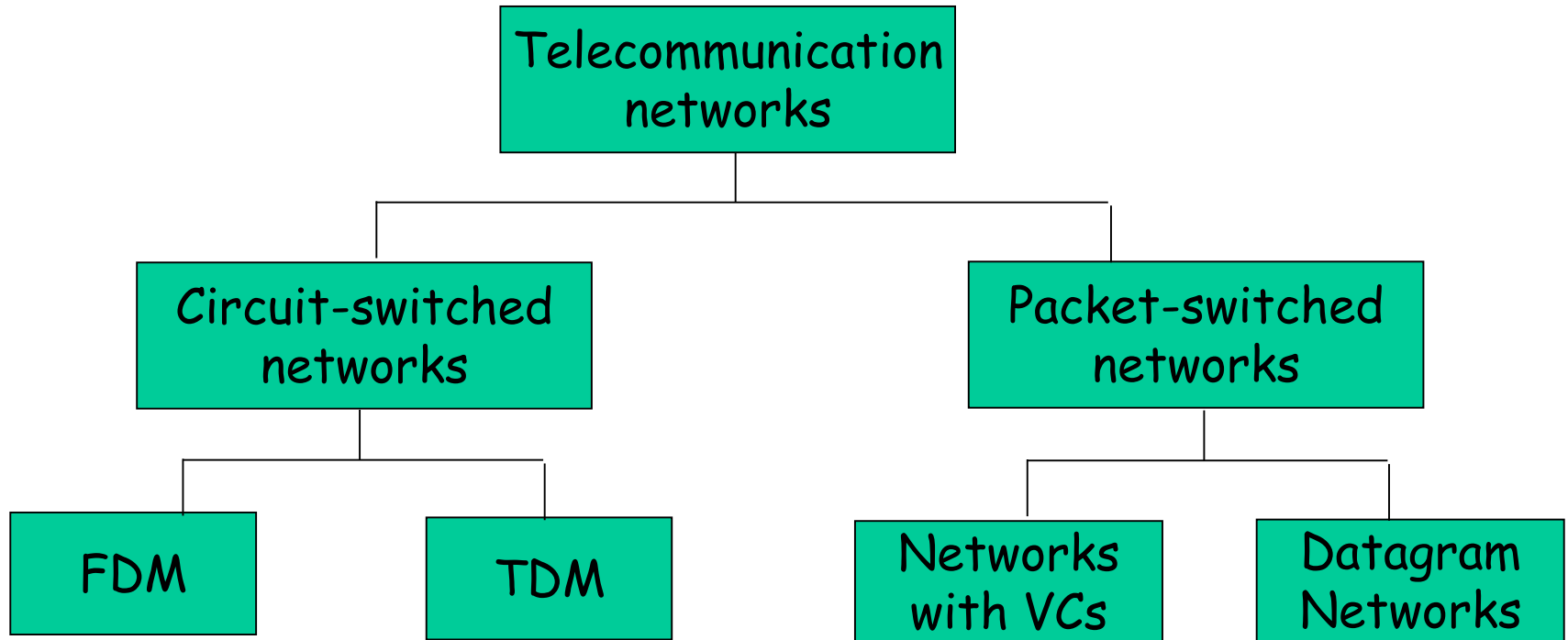
- *destination address* in packet determines next hop
- routes may change during session (flexible?)
- no “per flow” state, hence more scalable

□ virtual circuit network:

- each packet carries tag (virtual circuit ID), tag determines next hop
- fixed path determined at *call setup time*
- path is **not** a dedicated path as in circuit switched (i.e., store & forward of packets)
- *routers maintain per-call state*

□ datagram networks need per packet routing.

Network Taxonomy



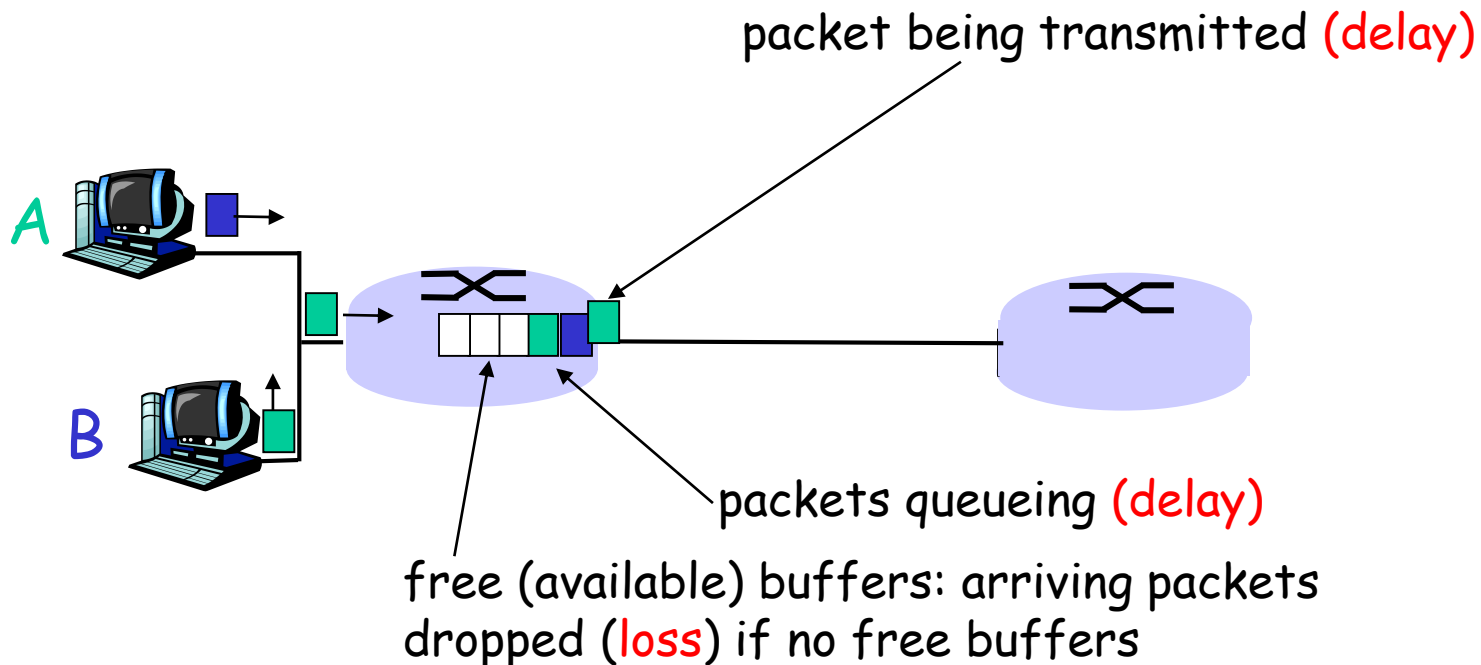
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How do loss and delay occur?

packets *queue* in router buffers

- ❑ packet arrival rate to link exceeds output link capacity
- ❑ packets queue, wait for turn
- ❑ if queue is full, arriving packets dropped (Drop-Tail)



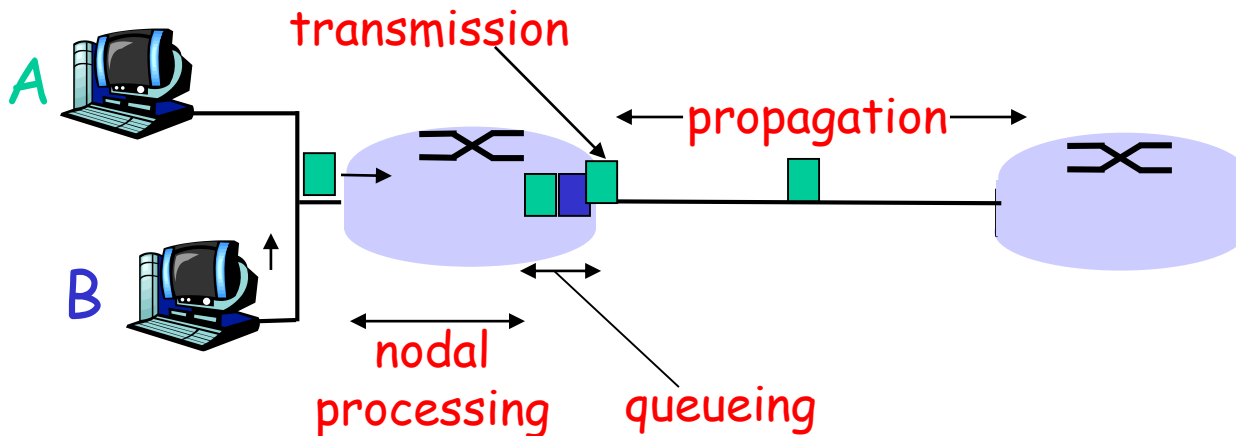
Four sources of packet delay

□ 1. nodal processing:

- check bit errors
- determine output link

□ 2. queueing

- time waiting at output link for transmission
- depends on congestion level of router



Delay in packet-switched networks

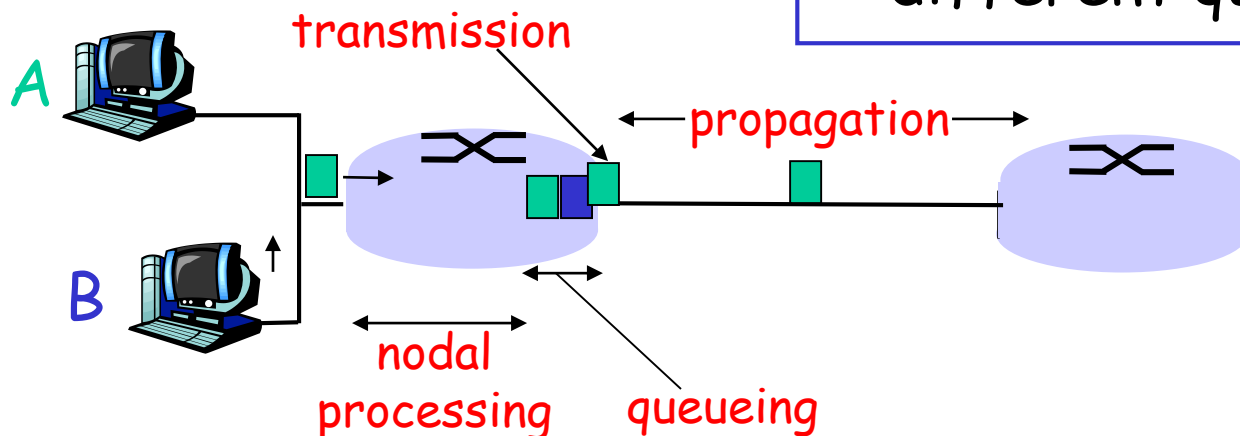
3. Transmission delay:

- R = link bandwidth (bps)
- L = packet length (bits)
- time to send bits into link = L/R

4. Propagation delay:

- d = length of physical link
- s = propagation speed in medium ($\sim 2 \times 10^8$ m/sec)
- propagation delay = d/s

Note: s and R are very different quantities!



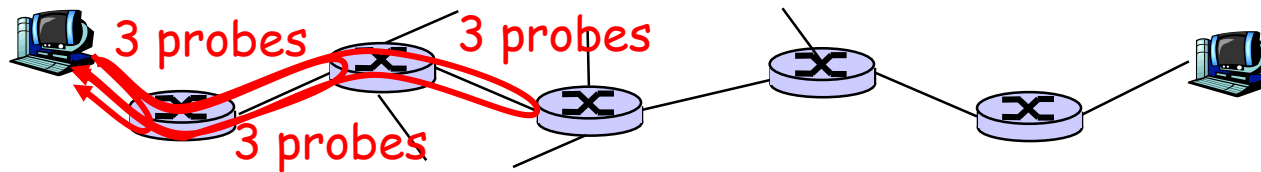
Nodal delay

$$d_{\text{nodal}} = d_{\text{proc}} + d_{\text{queue}} + d_{\text{trans}} + d_{\text{prop}}$$

- ❑ d_{proc} = processing delay
 - typically a few microsecs or less
- ❑ d_{queue} = queuing delay
 - depends on congestion
- ❑ d_{trans} = transmission delay
 - $= L/R$, significant for low-speed links
- ❑ d_{prop} = propagation delay
 - a few microsecs to hundreds of msecs

"Real" Internet delays and routes

- ❑ What do "real" Internet delay & loss look like?
- ❑ Traceroute program: provides delay measurement from source to router along end-end Internet path towards destination. For all i :
 - sends three packets that will reach router i on path towards destination
 - router i will return packets to sender
 - sender times interval between transmission and reply.

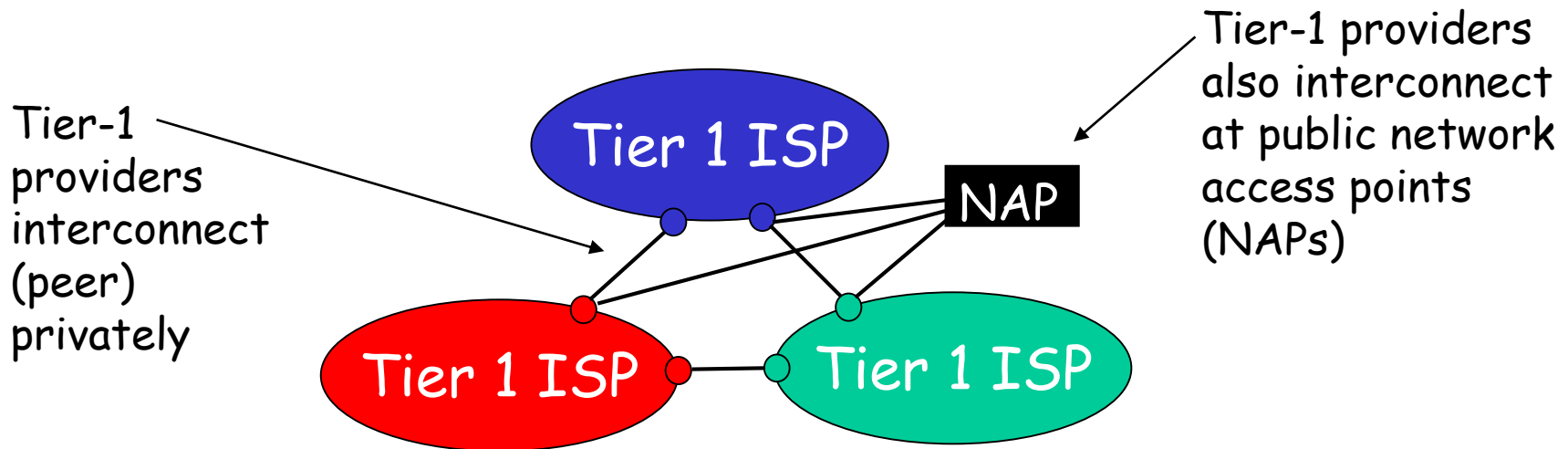


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Internet structure: network of networks

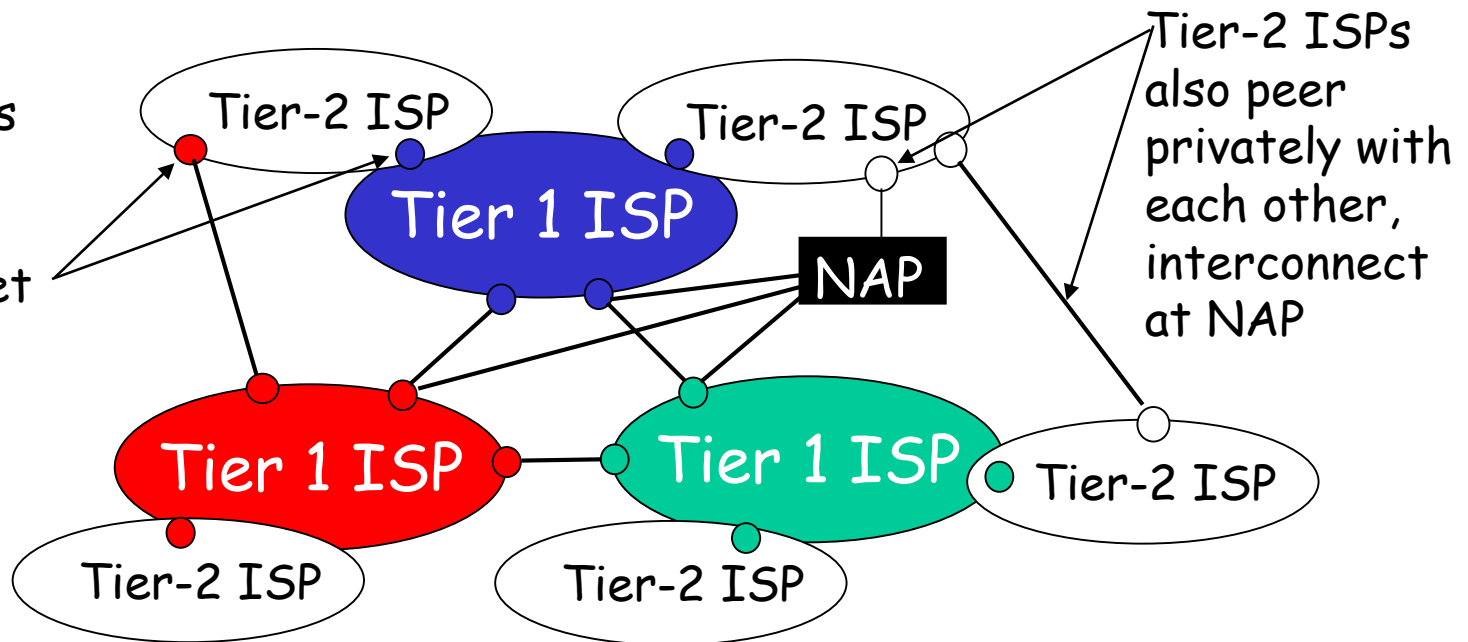
- ❑ roughly hierarchical
- ❑ **at center: "tier-1" ISPs** (e.g., UUNet, BBN/Genuity, Sprint, AT&T), national/international coverage
 - treat each other as equals



Internet structure: network of networks

- "Tier-2" ISPs: smaller (often regional) ISPs
 - Connect to one or more tier-1 ISPs, possibly other tier-2 ISPs

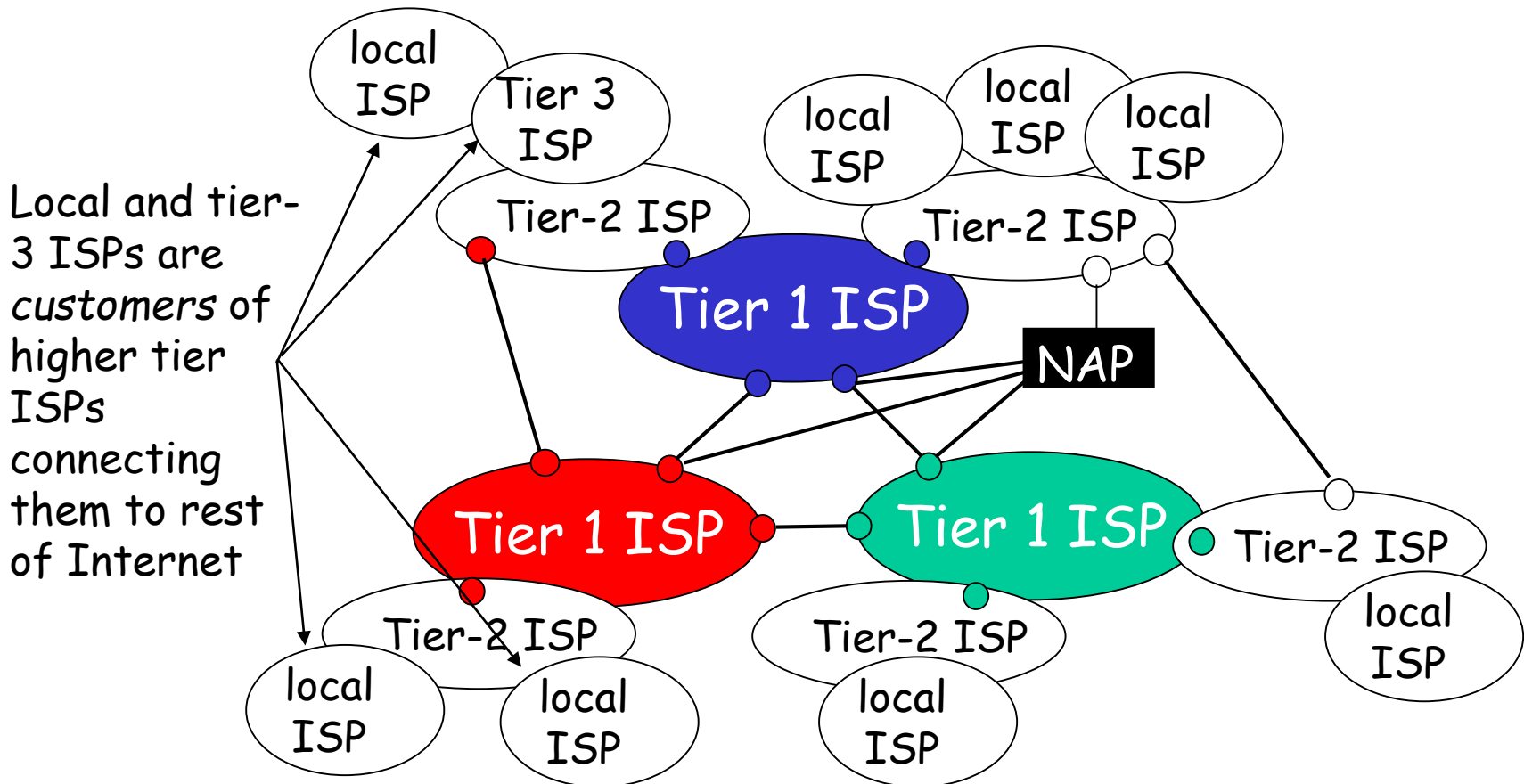
Tier-2 ISP pays tier-1 ISP for connectivity to rest of Internet
□ tier-2 ISP is customer of tier-1 provider



Internet structure: network of networks

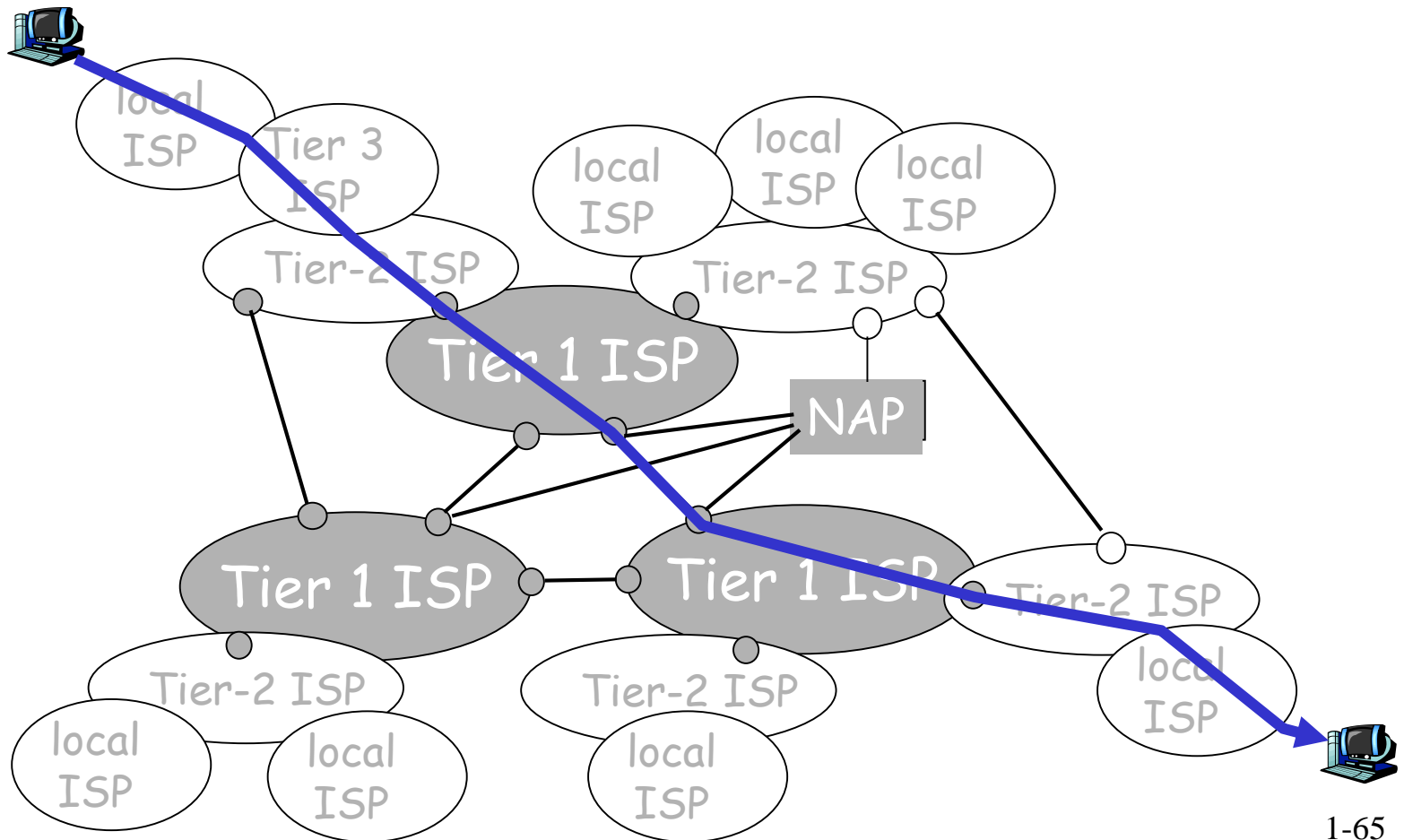
□ "Tier-3" ISPs and local ISPs

- last hop ("access") network (closest to end systems)



Internet structure: network of networks

- a packet passes through many networks!



Introduction: Summary

Covered a "ton" of material!

- ❑ Internet overview
- ❑ what's a protocol?
- ❑ network edge, core, access network
 - packet-switching versus circuit-switching
- ❑ Internet/ISP structure
- ❑ performance: loss, delay
- ❑ layering and service models
- ❑ history (which you will be reading on your own)

You now have:

- ❑ context, overview, "feel" of networking
- ❑ more depth, detail to follow!