CS 441 Discrete Mathematics for CS
Lecture 25

Graphs

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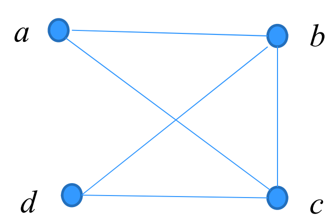
Definition of a graph

- **Definition:** A *graph* $G = (V, E)$ consists of a nonempty set V of *vertices* (or *nodes*) and a set E of *edges.* Each edge has either one or two vertices associated with it, called its *endpoints*. An edge is said to *connect* its endpoints.

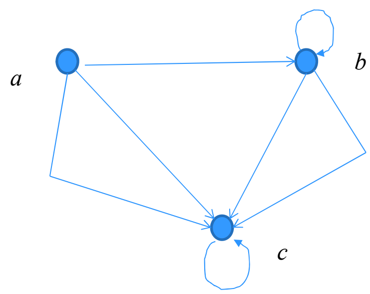
- **Example:**

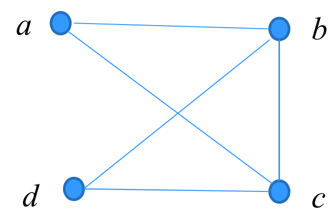


Graphs: basics

Basic types of graphs:

- **Directed graphs**
- **Undirected graphs**





Terminology

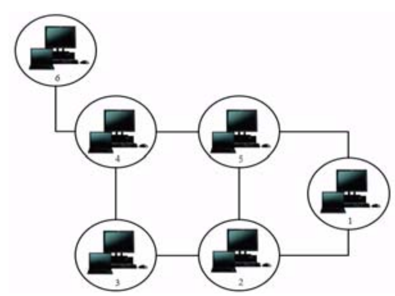
- In a *simple graph* each edge connects two different vertices and no two edges connect the same pair of vertices.
- *Multigraphs* may have multiple edges connecting the same two vertices. When m different edges connect the vertices u and v, we say that $\{u,v\}$ is an edge of *multiplicity* m.
- An edge that connects a vertex to itself is called a *loop*.
- A *pseudograph* may include loops, as well as multiple edges connecting the same pair of vertices.

Graphs

- **Graphs and graph theory can be used to model:**
  - Computer networks
  - Social networks
  - Communications networks
  - Information networks
  - Software design
  - Transportation networks
  - Biological networks

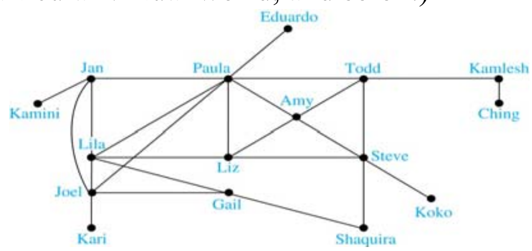
Graphs

- **Computer networks:**
  - **Nodes – computers**
  - **Edges - connections**



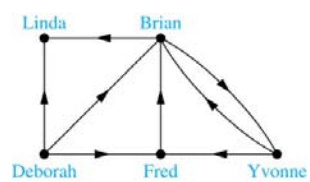
Graph models

- Graphs can be used to model social structures based on different kinds of relationships between people or groups.
- ***Social network***, vertices represent individuals or organizations and edges represent relationships between them.
- Useful graph models of social networks include:
  – ***friendship graphs*** - undirected graphs where two people are connected if they are friends (in the real world, on Facebook, or in a particular virtual world, and so on.)



Graph models

- Useful graph models of social networks include:
  – *influence graphs* - directed graphs where there is an edge from one person to another if the first person can influence the second person



  – *collaboration graphs* - undirected graphs where two people are connected if they collaborate in a specific way

Collaboration graphs

- The ***Hollywood graph*** models the collaboration of actors in films.
  - We represent actors by vertices and we connect two vertices if the actors they represent have appeared in the same movie.
  - Kevin Bacon numbers.
- An ***academic collaboration graph*** models the collaboration of researchers who have jointly written a paper in a particular subject.
  - We represent researchers in a particular academic discipline using vertices.
  - We connect the vertices representing two researchers in this discipline if they are coauthors of a paper.
  - *Erdős number*.

Information graphs

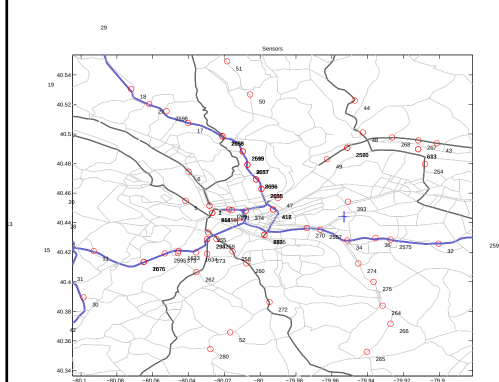
- Graphs can be used to model different types of networks that link different types of information.
- In a ***web graph***, web pages are represented by vertices and links are represented by directed edges.
  - A web graph models the web at a particular time.
- In a ***citation network***:
  - Research papers in a particular discipline are represented by vertices.
  - When a paper cites a second paper as a reference, there is an edge from the vertex representing this paper to the vertex representing the second paper.

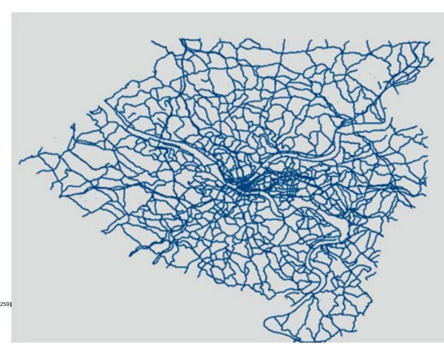
Transportation graphs

- Graph models are extensively used in the study of transportation networks.
- ***Airline networks*** modeled using directed multigraphs:
  - airports are represented by vertices
  - each flight is represented by a directed edge from the vertex representing the departure airport to the vertex representing the destination airport
- ***Road networks*** can be modeled using graphs where
  - vertices represent intersections and edges represent roads.
  - undirected edges represent two-way roads and directed edges represent one-way roads.

Transportation graphs

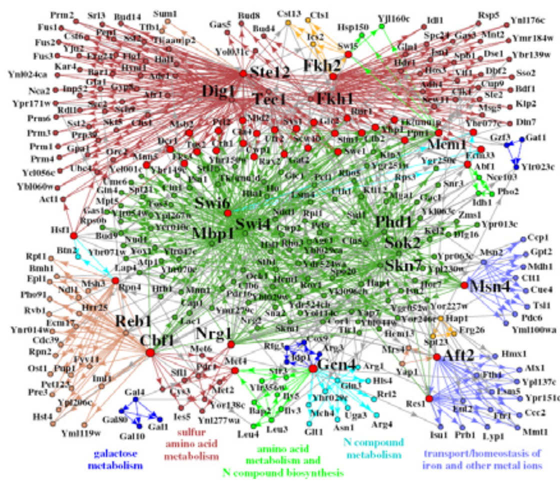
- Graph models are extensively used in the study of transportation networks.





Graphs

- **Biological networks:**



Graph characteristics: Undirected graphs

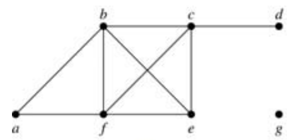
Definition 1. Two vertices *u*, *v* in an undirected graph *G* are called ***adjacent*** **(or *neighbors*)** in *G* if there is an edge *e* between *u* and *v*. Such an edge *e* is called *incident with* the vertices *u* and *v* and *e* is said to *connect* *u* and *v*.

Definition 2. The set of all neighbors of a vertex *v* of *G* = (*V*, *E*), denoted by *N*(*v*), is called **the *neighborhood* of *v***. If *A* is a subset of *V*, we denote by *N*(*A*) the set of all vertices in *G* that are adjacent to at least one vertex in *A*. So,

Definition 3. The ***degree of a vertex in a undirected graph*** is the number of edges incident with it, except that a loop at a vertex contributes two to the degree of that vertex. The degree of the vertex *v* is denoted by deg(*v*).

Undirected graphs

Example: What are the degrees and neighborhoods of the vertices in the graphs G?



G

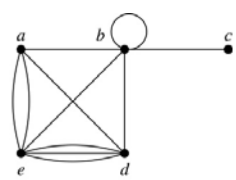
Solution:

G: $\deg(a) = 2$, $\deg(b) = \deg(c) = \deg(f) = 4$, $\deg(d) = 1$, $\deg(e) = 3$, $\deg(g) = 0$.

$N(a) = \{b, f\}$, $N(b) = \{a, c, e, f\}$, $N(c) = \{b, d, e, f\}$,
$N(d) = \{c\}$, $N(e) = \{b, c, f\}$, $N(f) = \{a, b, c, e\}$, $N(g) = \varnothing$.

Undirected graphs

Example: What are the degrees and neighborhoods of the vertices in the graphs H?



H

Solution:

H: $\deg(a) = 4$, $\deg(b) = \deg(e) = 6$, $\deg(c) = 1$, $\deg(d) = 5$.

$N(a) = \{b, d, e\}$, $N(b) = \{a, b, c, d, e\}$, $N(c) = \{b\}$,
$N(d) = \{a, b, e\}$, $N(e) = \{a, b, d\}$.

Undirected graphs

Theorem 1 (*Handshaking Theorem*): If $G = (V,E)$ is an undirected graph with m edges, then

$$2m = \sum_{v \in V} \deg(v)$$

Proof:

Each edge contributes twice to the degree count of all vertices. Hence, both the left-hand and right-hand sides of this equation equal twice the number of edges.

Think about the graph where vertices represent the people at a party and an edge connects two people who have shaken hands.

Undirected graphs

Theorem 2: An undirected graph has an even number of vertices of odd degree.

***Proof*:** Let V_1 be the vertices of even degree and V_2 be the vertices of odd degree in an undirected graph $G = (V, E)$ with m edges. Then

$$2m = \sum_{v \in V} \deg(v) = \sum_{v \in V_1} \deg(v) + \sum_{v \in V_2} \deg(v).$$

must be even since $\deg(v)$ is even for each $v \in V_1$

This sum must be even because $2m$ is even and the sum of the degrees of the vertices of even degrees is also even. Because this is the sum of the degrees of all vertices of odd degree in the graph, there must be an even number of such vertices.

Directed graphs

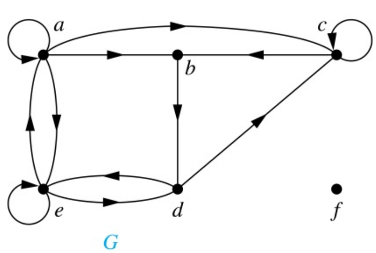
Definition: An ***directed graph*** $G = (V, E)$ consists of *V,* a nonempty set of *vertices* (or *nodes*), and *E,* a set of *directed edges* or *arcs.* Each edge is an ordered pair of vertices. The directed edge (u,v) is said to start at *u* and end at *v*.

Definition: Let (u,v) be an edge in *G*. Then *u* is the *initial vertex* of this edge and is *adjacent to v* and *v* is the *terminal* (or *end*) *vertex* of this edge and is *adjacent from u*. The initial and terminal vertices of a loop are the same.

Directed graphs

Definition: The *in-degree of a vertex v*, denoted $deg^{-}(v)$, is the number of edges which terminate at *v*. The *out-degree of v*, denoted $deg^{+}(v)$, is the number of edges with *v* as their initial vertex. Note that a loop at a vertex contributes 1 to both the in-degree and the out-degree of the vertex.

Example: Assume graph *G:*



What are in-degrees of vertices: ?

$\deg^{-}(a) = 2$, $\deg^{-}(b) = 2$, $\deg^{-}(c) = 3$, $\deg^{-}(d) = 2$, $\deg^{-}(e) = 3$, $\deg^{-}(f) = 0$.

Graphs: basics

Definition: The *in-degree of a vertex v*, denoted *deg*$^-$(v), is the number of edges which terminate at v. The *out-degree of v*, denoted *deg*$^+$(v), is the number of edges with v as their initial vertex. Note that a loop at a vertex contributes 1 to both the in-degree and the out-degree of the vertex.

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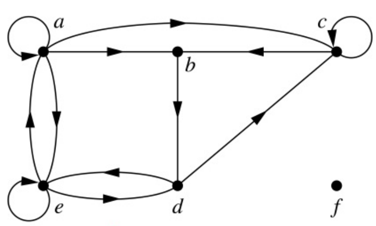


What are out-degrees of vertices: ?

$\deg^+(a) = 4, \deg^+(b) = 1, \deg^+(c) = 2,$
$\deg^+(d) = 2, \deg^+(e) = 3, \deg^+(f) = 0.$

Directed graphs

Theorem: Let $G = (V, E)$ be a graph with directed edges. Then:

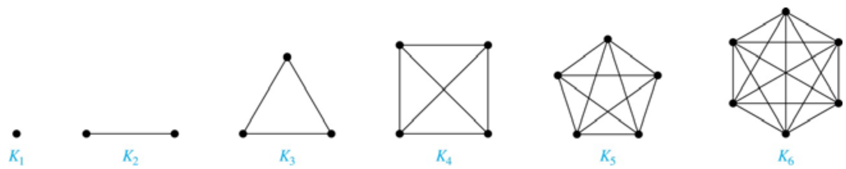
$$|E| = \sum_{v \in V} deg^-(v) = \sum_{v \in V} deg^+(v).$$

Proof:

The first sum counts the number of outgoing edges over all vertices and the second sum counts the number of incoming edges over all vertices. It follows that both sums equal the number of edges in the graph.

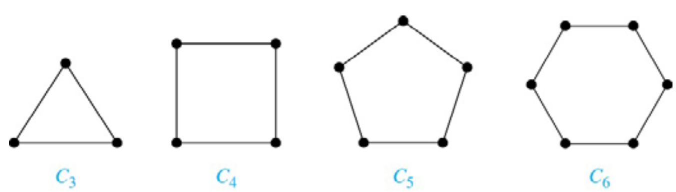
Complete graphs

A *complete graph on n vertices*, denoted by K_n, is the simple graph that contains exactly one edge between each pair of distinct vertices.



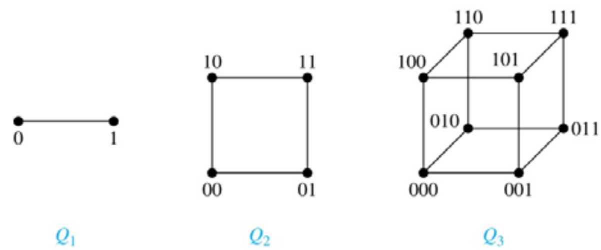
A cycle

A *cycle* C_n for $n \geq 3$ consists of n vertices $v_1, v_2, \cdots, v_n$, and edges $\{v_1, v_2\}, \{v_2, v_3\}, \cdots, \{v_{n-1}, v_n\}, \{v_n, v_1\}$.



N-dimensional hypercube

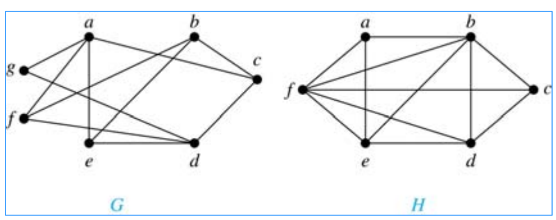
An *n-dimensional hypercube*, or *n-cube,* Q_n, is a graph with 2^n vertices representing all bit strings of length *n*, where there is an edge between two vertices that differ in exactly one bit position.



Bipartite graphs

Definition: A simple graph *G* is **bipartite** if *V* can be partitioned into two disjoint subsets V_1 and V_2 such that every edge connects a vertex in V_1 and a vertex in V_2. In other words, there are no edges which connect two vertices in V_1 or in V_2.

Note: An equivalent definition of a bipartite graph is a graph where it is possible to color the vertices red or blue so that no two adjacent vertices are the same color.



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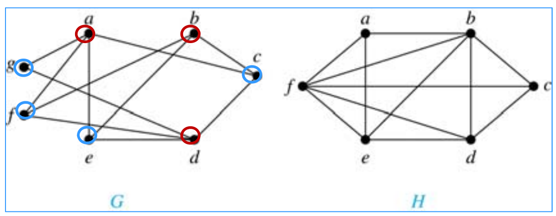
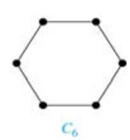


Bipartite graphs

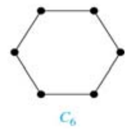
Example: Show that C_6 is bipartite.



Solution:

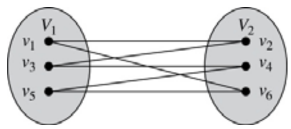
Bipartite graphs

Example: Show that C_6 is bipartite.



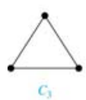
Solution:

- We can partition the vertex set into $V_1 = \{v_1, v_3, v_5\}$ and $V_2 = \{v_2, v_4, v_6\}$ so that every edge of C_6 connects a vertex in V_1 and V_2 .



Bipartite graphs

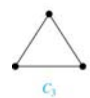
Example: Show that C_3 is not bipartite.



Solution:

Bipartite graphs

Example: Show that C_3 is not bipartite.



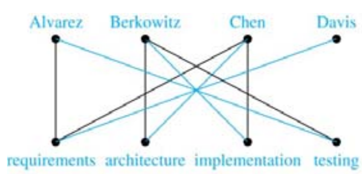
Solution:

If we divide the vertex set of C_3 into two nonempty sets, one of the two must contain two vertices. But in C_3 every vertex is connected to every other vertex. Therefore, the two vertices in the same partition are connected. Hence, C_3 is not bipartite.

Bipartite graphs and matching

Bipartite graphs are used to model applications that involve **matching** the elements of one set to elements in another, for example:

Example: *Job assignments* - vertices represent the jobs and the employees, edges link employees with those jobs they have been trained to do. A common goal is to match jobs to employees so that the most jobs are done.

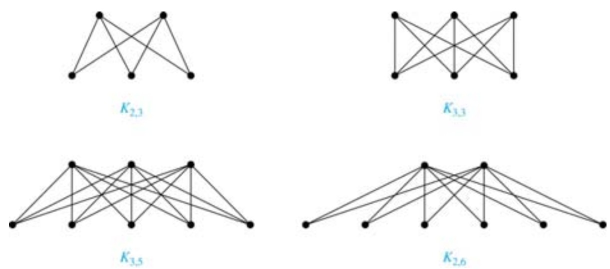


(a)

Complete bipartite graphs

Definition: A *complete bipartite graph* $K_{m,n}$ is a graph that has its vertex set partitioned into two subsets V_1 of size m and V_2 of size n such that there is an edge from every vertex in V_1 to every vertex in V_2.

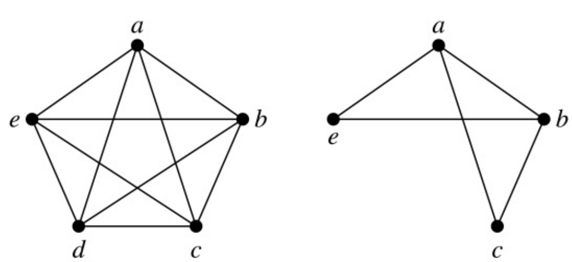
Example: We display four complete bipartite graphs here.



Subgraphs

Definition: A ***subgraph of a graph*** $G = (V,E)$ is a graph (W,F), where $W \subset V$ and $F \subset E$. A subgraph H of G is a proper subgraph of G if $H \neq G$.

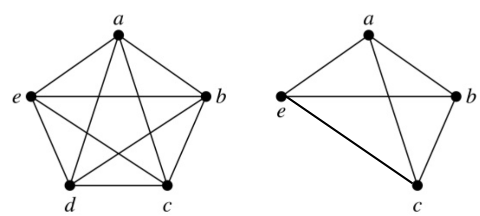
Example: K_5 and one of its subgraphs.



Subgraphs

Definition: Let $G = (V, E)$ be a simple graph. The *subgraph induced* by a subset W of the vertex set V is the graph (W,F), where the edge set F contains an edge in E if and only if both endpoints are in W.

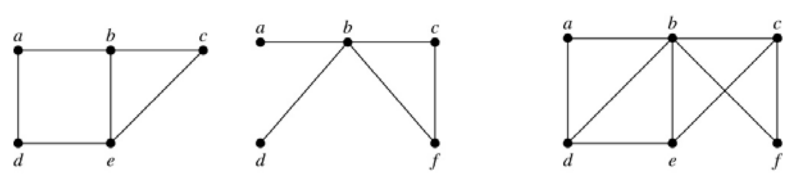
Example: K_5 and the subgraph induced by $W = \{a,b,c,e\}$.



Union of the graphs

Definition: The *union* of two simple graphs $G_1 = (V_1, E_1)$ and $G_2 = (V_2, E_2)$ is the simple graph with vertex set $V_1 \cup V_2$ and edge set $E_1 \cup E_2$. The union of G_1 and G_2 is denoted by $G_1 \cup G_2$.

Example:



G_1 G_2 $G_1 \cup G_2$

(a) (b)